

Dynamic Program Analyses and Their Security Applications

Xiangyu Zhang



What is Dynamic Program Analysis

- Dynamic analysis analyzes program execution
 - Examples
 - Identify the sequence of instructions that get executed
 - Identify the set of values that a variable holds during its lifetime
 - Identify the set of memory addresses/files that are accessed
 - Static analysis analyzes programs without executing them
 - Identify loops in the program
 - Generate control flow graph
 - Identify code clones
 - Static tainting

Applications of Dynamic Analysis

- Software engineering
 - Debugging: what is the root cause that leads to an observed program failure?
 - Software optimization
 - Software validation
- Security
 - Anomaly detection
 - Forensic analysis
 - Expose hidden malicious logic

Outline

Temporal Dynamic Analysis:

analyze program execution history

1. Attack investigation

2. Application vetting

Spatial Dynamic Analysis:

analyze a snapshot of program execution state

3. Memory forensics

Dynamic Analyses

Outline

Temporal Dynamic Analysis:

analyze program execution history

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Spatial Dynamic Analysis:

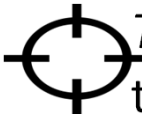
analyze a snapshot of program execution state

3. Memory forensics

Dynamic Analyses

Cyberattacks are becoming more sophisticated

- **Advanced Persistent Threat (APT)**

 *Targeted:* Targets specific organizations to exfiltrate information or disrupt the systems.

Infrastructure
(Nuclear plants)

Business
(Target® Data Breach)

Government
(OPM: Office of Personnel
Management)

Politics
(DNC email hack)



Multiple stages of APTs



1. *Reconnaissance*: Learn the target organization



2. *Infiltration*: Enter into the victim via social-engineering (e.g., phishing emails) or vulnerabilities



3. *Discovery and capture*: Stay *low* and operate *slowly* to avoid detection while discovering critical machines and/or information.



4. *Exfiltration/Disruption*: Send the captured secret information to attackers or destroy the systems

Combatting APTs is challenging



3. *Discovery and capture:* Stay *low* and operate *slowly* to avoid detection while discovering critical machines and/or information.

(Whitelisted) benign built-in software

APT actors often leverage *benign built-in software* (e.g., *web-browsers and email clients* that are *already whitelisted*) to avoid detection

Low and slow (Stealthy)

Incidents are often detected after a *few months*.

Audit Logging

- Audit logging records system level events during system execution
 - For attack investigation: forensic analysis
 - Identify the source of an attack
 - Understand the damage to a victim system



Audit Logging

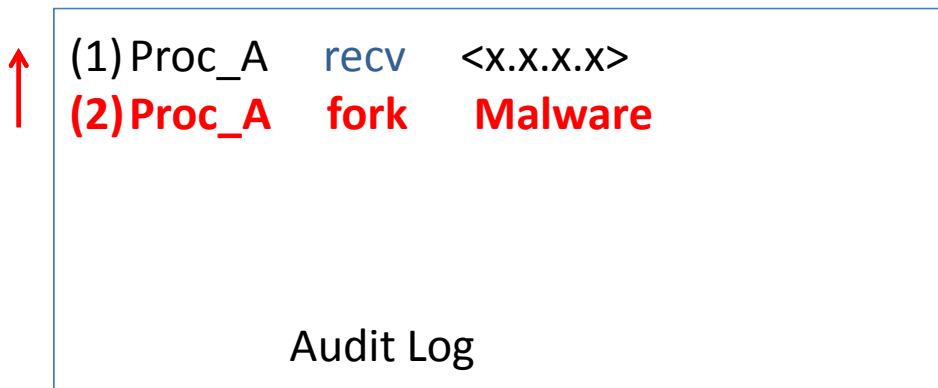
- Analyze audit logs to generate a causal graph

```
(1) Proc_A  recv  <x.x.x.x>  
(2) Proc_A  fork  Malware
```

Audit Log

Audit Logging

- Analyze audit logs to generate a causal graph
 - Backward analysis** - identify the source of an attack



○ : Process

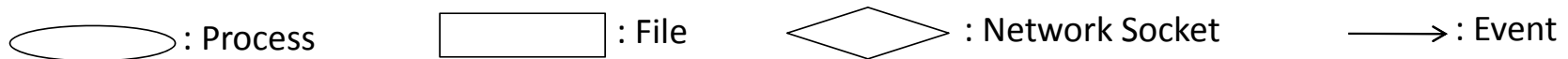
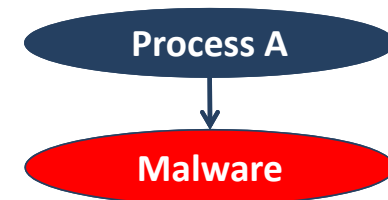
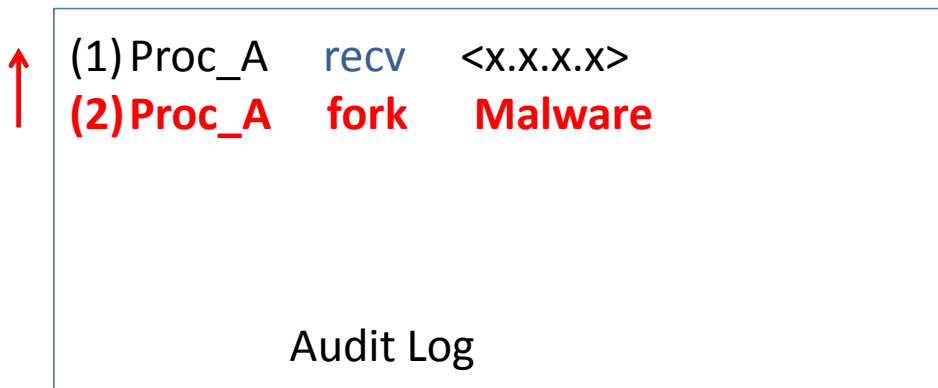
□ : File

◇ : Network Socket

→ : Event

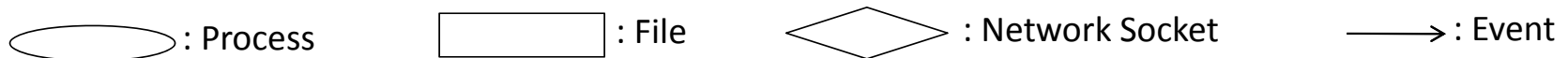
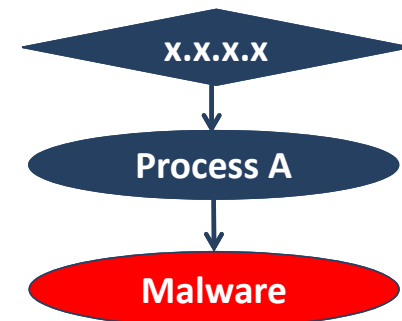
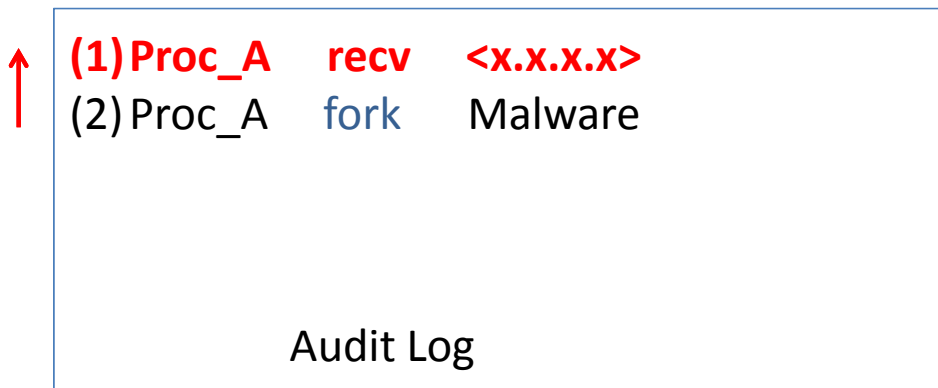
Audit Logging

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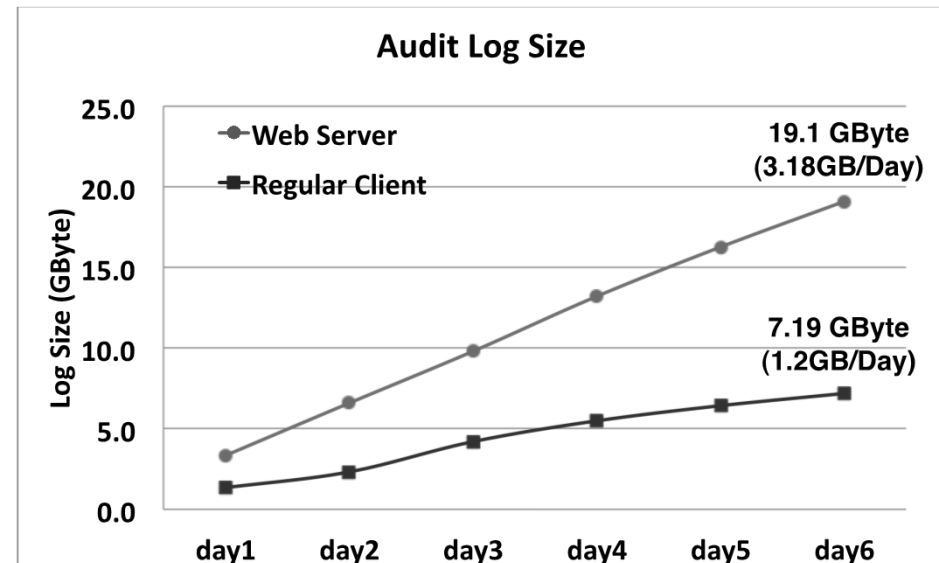
Audit Logging

- Analyze audit logs to generate a causal graph
 - Backward analysis** - identify the source of an attack



Limitations of Traditional Audit Logging

- Dependence explosion problem
 - Long running programs generate a lot of system level dependences during their lifetime
- High overhead
 - Linux Audit Framework: **~40%** run time slow down
 - Space overhead



Dependence Explosion - Example

- Social engineering attack by phishing e-mail



From: Chase Online [mailto:smrfs@chaseonline.com]
Sent: Wednesday, July 11, 2012 8:27 PM
Subject: Verification of Recent Activities Required



URGENT: Verification of Recent Activities Required
Your Chase Bank Account

Dear Customer:

As part of our ongoing effort to protect your account and our relationship, we monitor your account for possible fraudulent activity. **We need to confirm that you or someone authorized to use your account made the following sign in error attempt on your Chase Bank account:**

- 1) Sign in Error Attempt was noticed and registered at 70.43.95.130, Chantilly, Virginia United State on or around 2012-07-11 at 05:01AM.
- 2) Sign in Error Attempt was noticed and registered at 68.170.136.81, Commack, New York United State on or around 2012-07-11 at 8:30PM.
- 3) Sign in Error Attempt was noticed and registered at 74.11.185.43 Delray Beach, Florida United State on or around 2012-07-11 at 8:20PM.
- 4) Sign in Error Attempt was noticed and registered at 68.46.148.86, Egg Harbor Township, New Jersey, United States on or around 2012-07-11 at 6:39AM.

Please click on the link below to sign in correctly to re-activate your online banking access:

www.chase.com

Your satisfaction is important to us, and we appreciate your prompt attention to this matter. If you already had the opportunity to discuss this matter with us, please disregard this message.

Thank you for being our customer.

Sincerely,

A handwritten signature in black ink, appearing to read "Chris Palumbo".

Christopher J. Palumbo
Senior Vice President
Chase Fraud Prevention

- Social engin



[My Accounts](#) > [Download Activity](#)

Download Activity

[Help with this page](#)

Download account transactions — Download your transaction details into your Quicken®, QuickBooks® or Microsoft® Money software. Use the drop-down list to select your software format, then click "Download Activity." **Note:** To download transactions from the last 45 days, leave the beginning and ending date fields blank.

Attention! WaMu credit card customers: Before you attempt to download transactions into your Quicken, QuickBooks or Microsoft Money software, please [use these special instructions](#) to help you update your software on or after **March 9**.

*Required field

Download Information

Select account*

Select Account

Choose date range*

- ☒ All transactions available (Limited to 45 days)
☐ Specify a date range

Beginning date

(mm/dd/yyyy)

Ending date

(mm/dd/yyyy)

Select software format*

-- Select Download Type --

*Required field

Download Activity

Cancel



Download Activity

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Download account transactions — Download your transaction details into your Quicken®, QuickBooks® or Microsoft® Money software. Use the drop-down list to select your software format, then click "Download Activity." **Note:** To download transactions from the last 45 days, leave the beginning and ending date fields blank.

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-- Select Download Type --

*Required field

[Download Activity](#)

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Download Activity

Download account trans
into your Quicken®, QuickBooks® or I
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Attention! WaMu credit card custom
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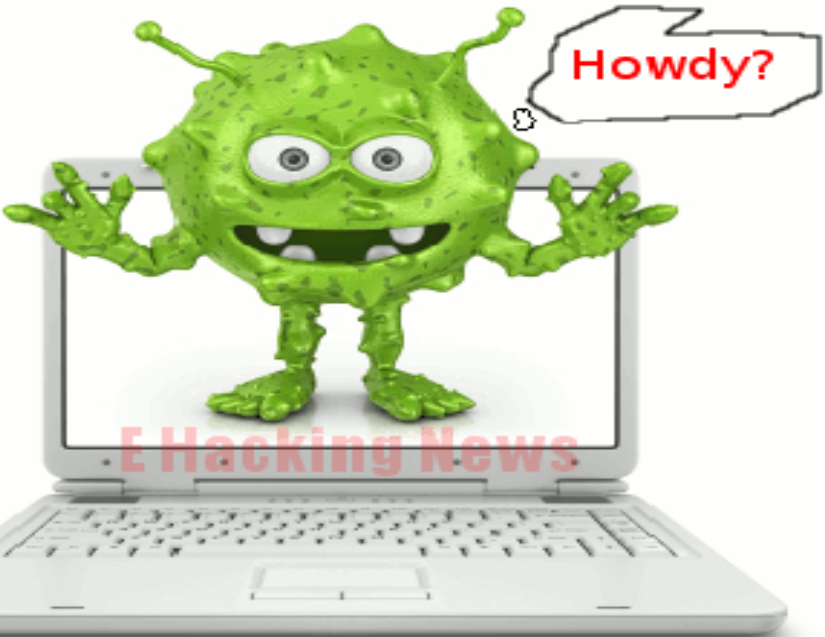
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[*Backward Analysis*]

Malware

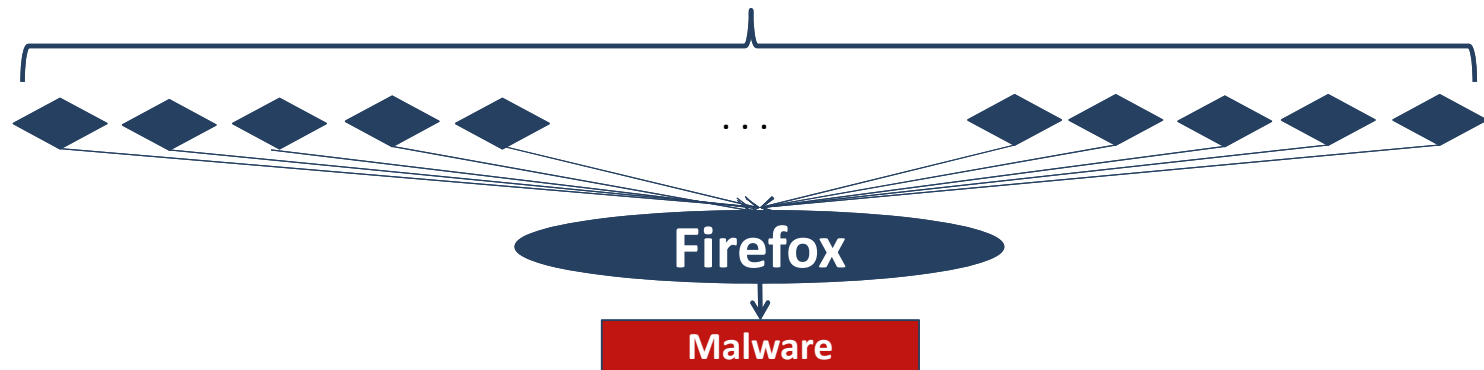
[*Backward Analysis*]



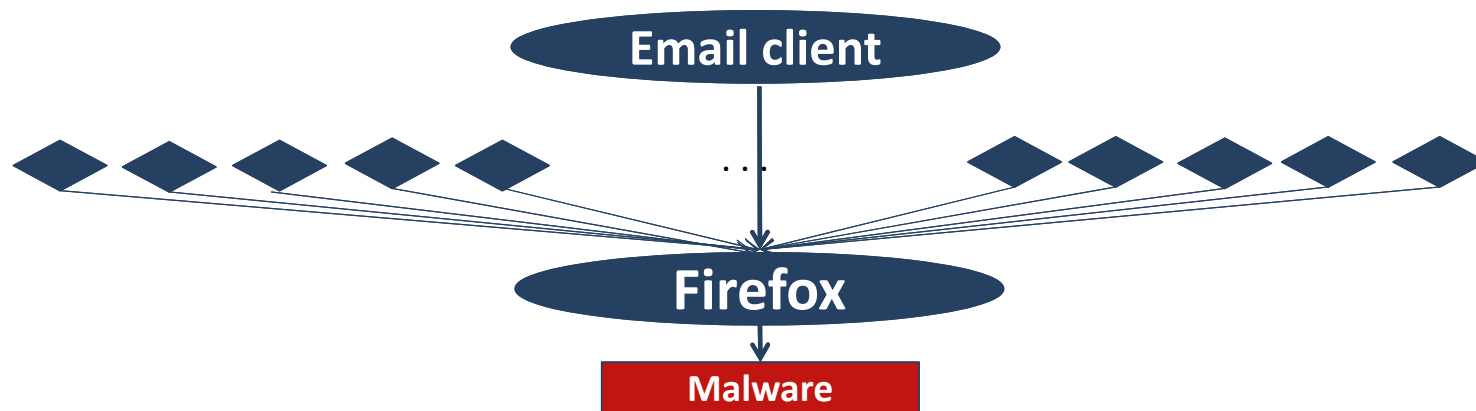
[*Backward Analysis*]

The user visited 11 web sites

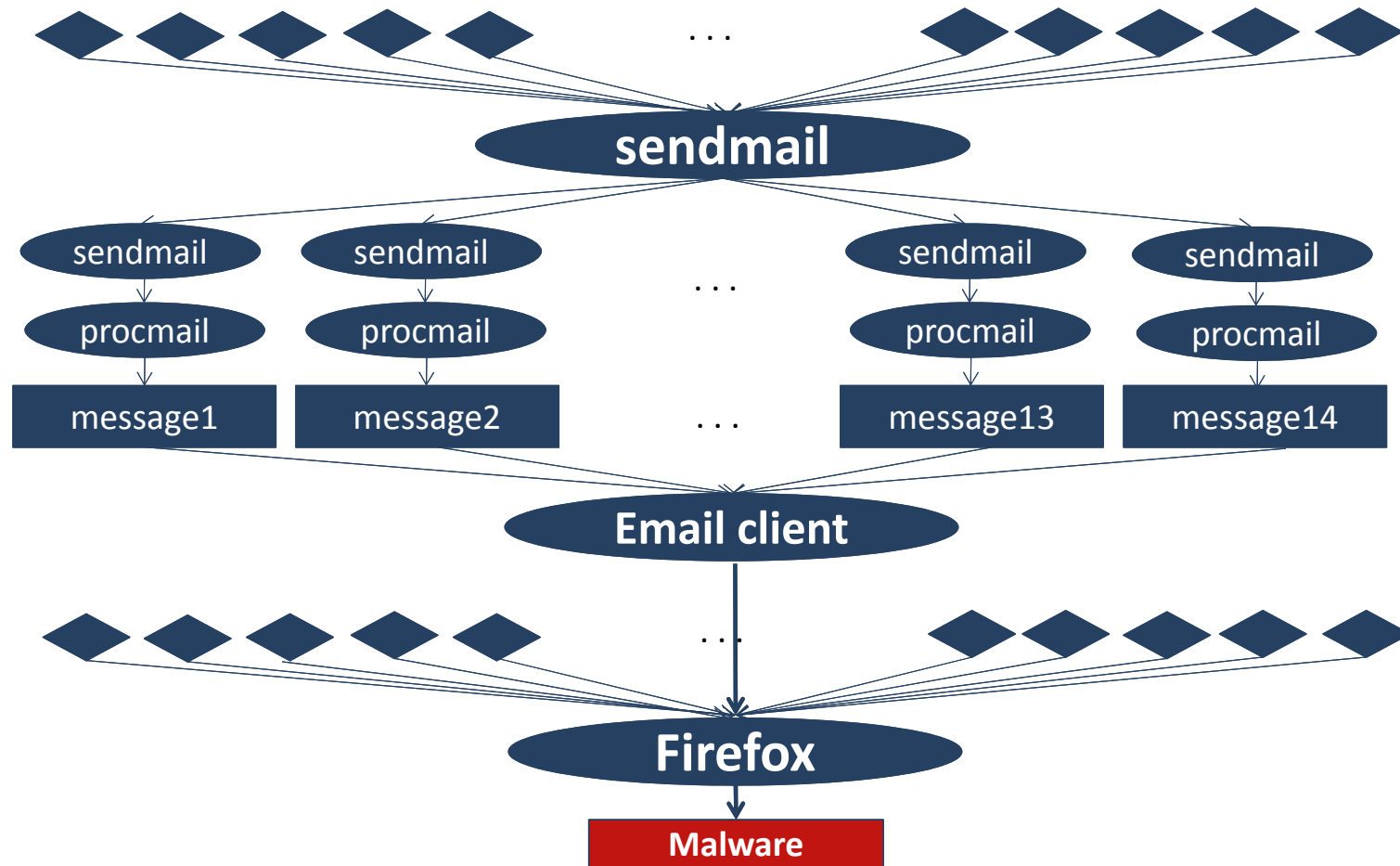
Dependence explosion!!
(229 IP Addresses)



[*Backward Analysis*]



[*Backward Analysis*]



[*Backward Analysis*]

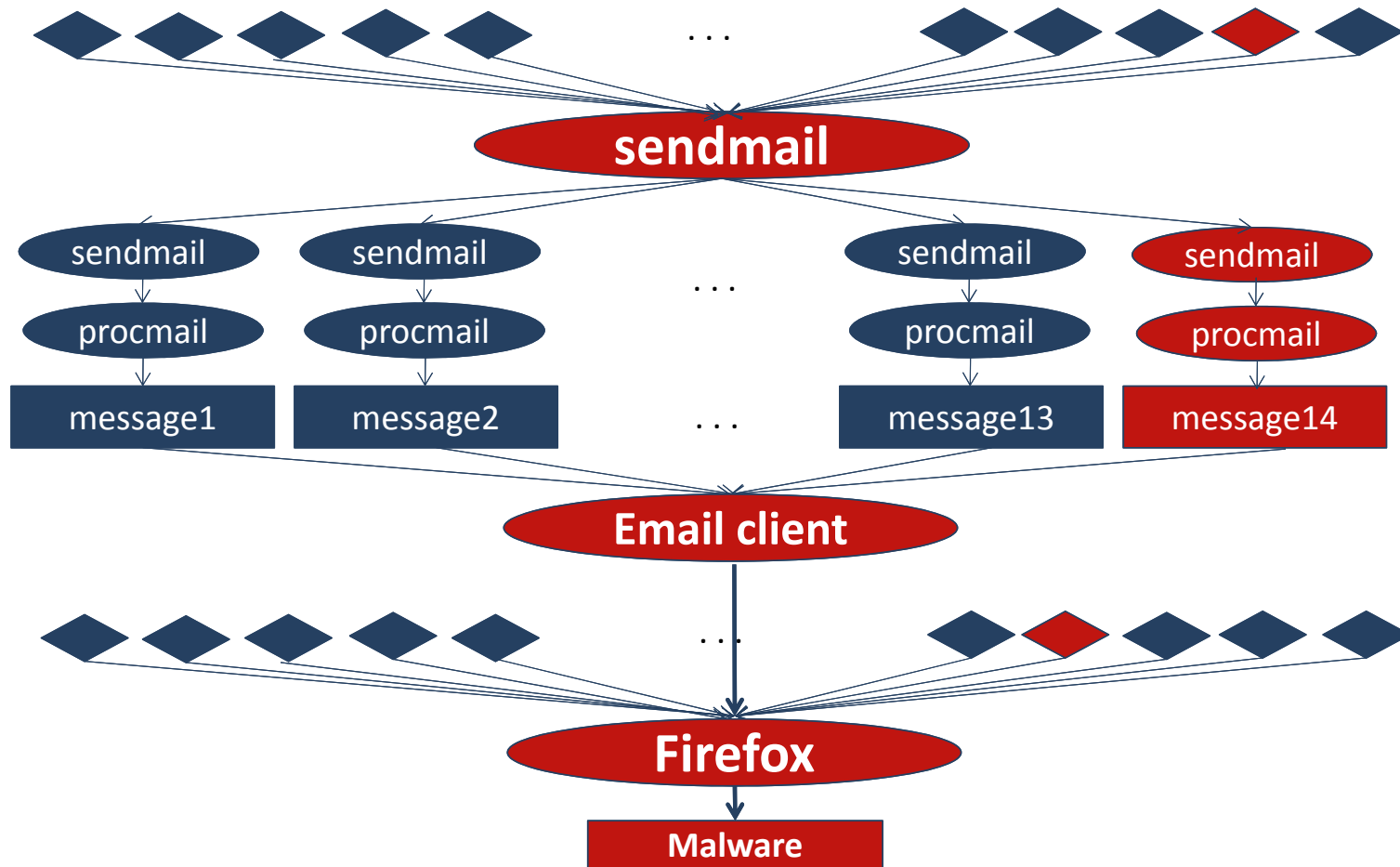
Dependence Explosion :
51 Processes, 15 Files,
251 Network addresses, 351 Edges



Malware

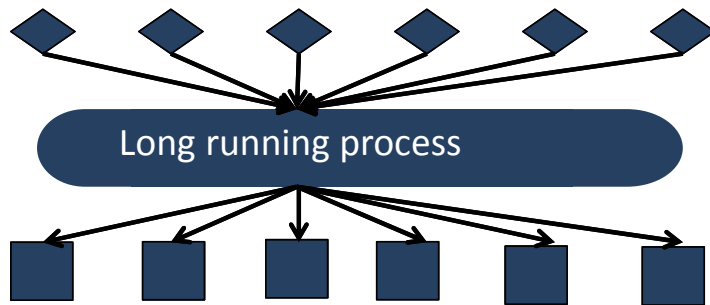
The diagram illustrates the flow of email data from a mail server to a malware-infected device. At the top, a central oval labeled "sendmail" receives input from multiple diamond-shaped nodes (representing email clients or devices) and sends output to several "procmail" ovals. Each "procmail" oval is connected to a "message" rectangle (e.g., "message1", "message2", "message13", "message14"). These messages are then sent to a central oval labeled "Email client". The "Email client" is connected to a "Firefox" oval, which in turn is connected to a "Malware" rectangle at the bottom. The diagram shows a clear path from the mail server to the malware-infected device.

[*Backward Analysis*]



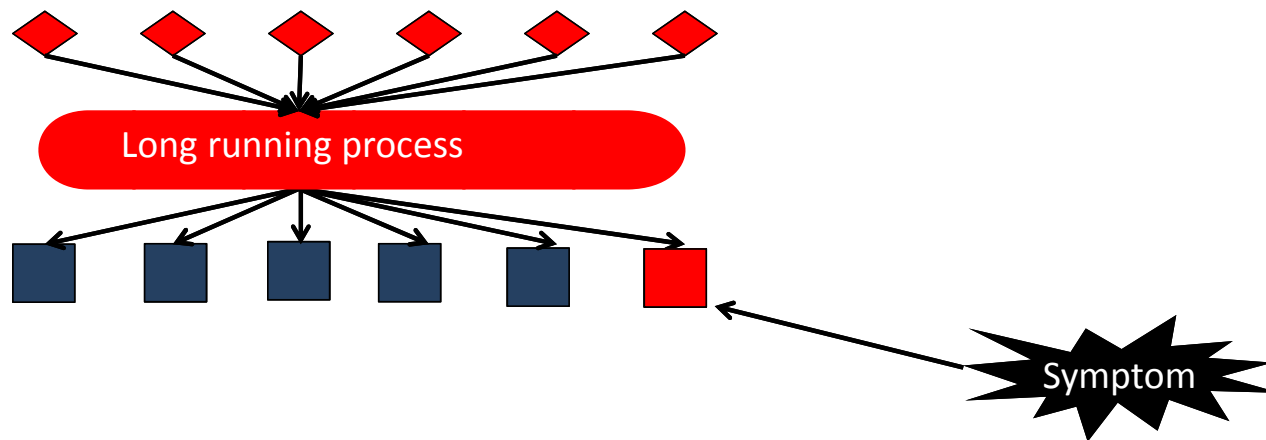
Dependence Explosion – Root Cause

- Caused by **long-running processes**
 - Receive many inputs and produces many outputs



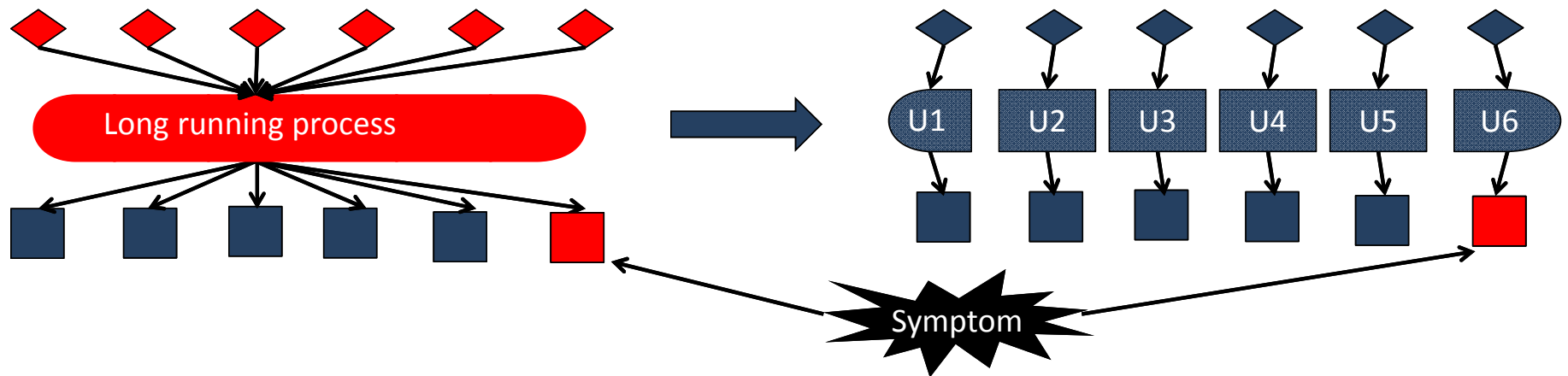
Dependence Explosion – Root Cause

- Caused by **long-running processes**
 - Receive many inputs and produces many outputs
 - Any output is potentially related to all preceding inputs



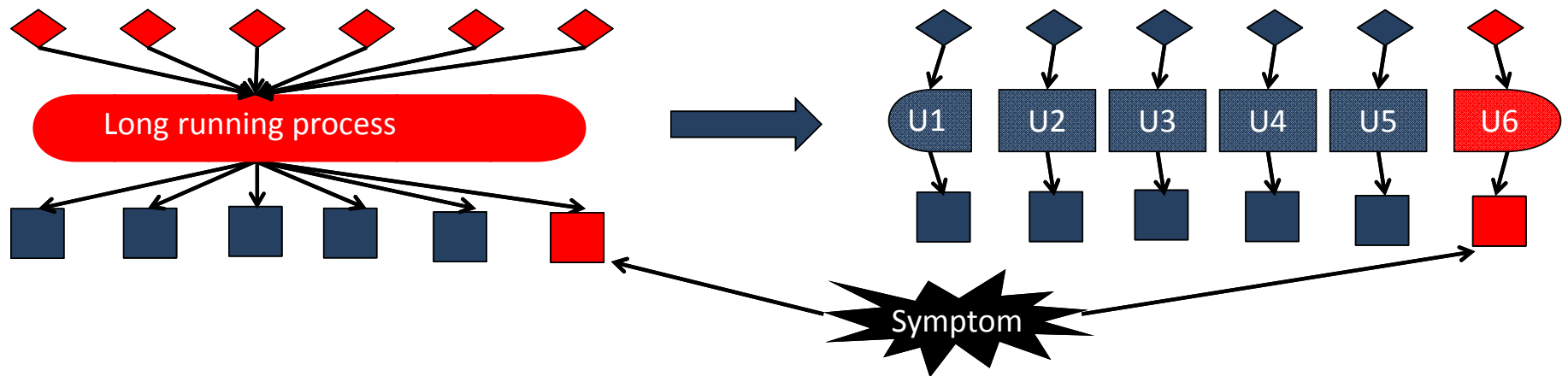
BEEP : Binary-based ExEcution Partition

- Finer-grained subject : Execution “UNIT”
 - Dynamically partition the execution of a process into **autonomous** execution segments



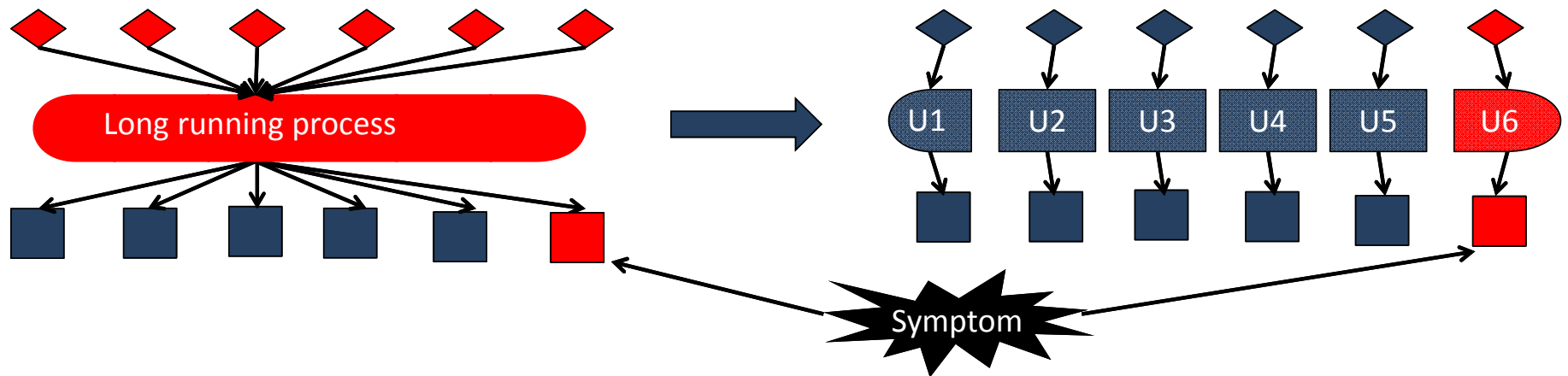
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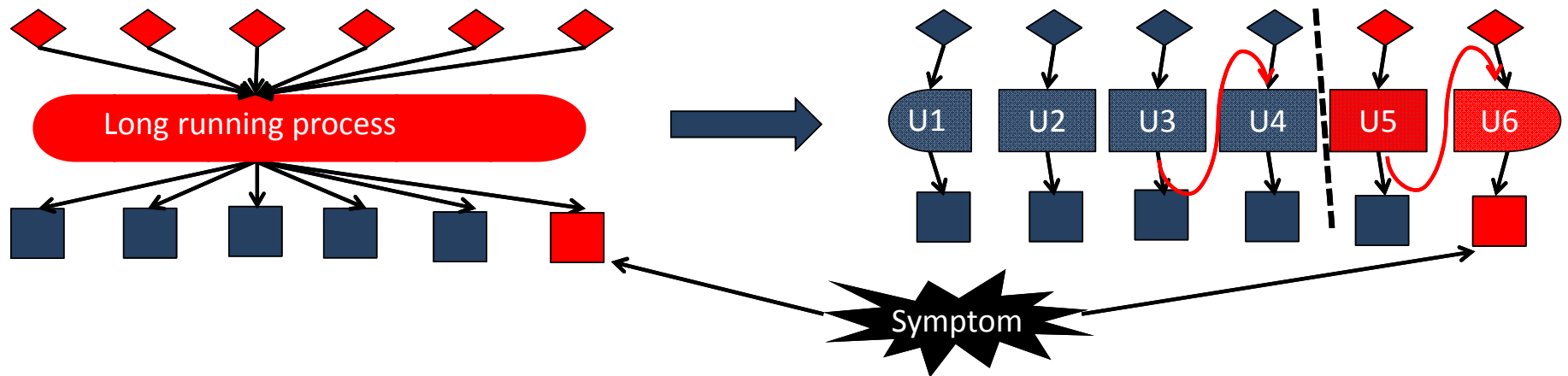
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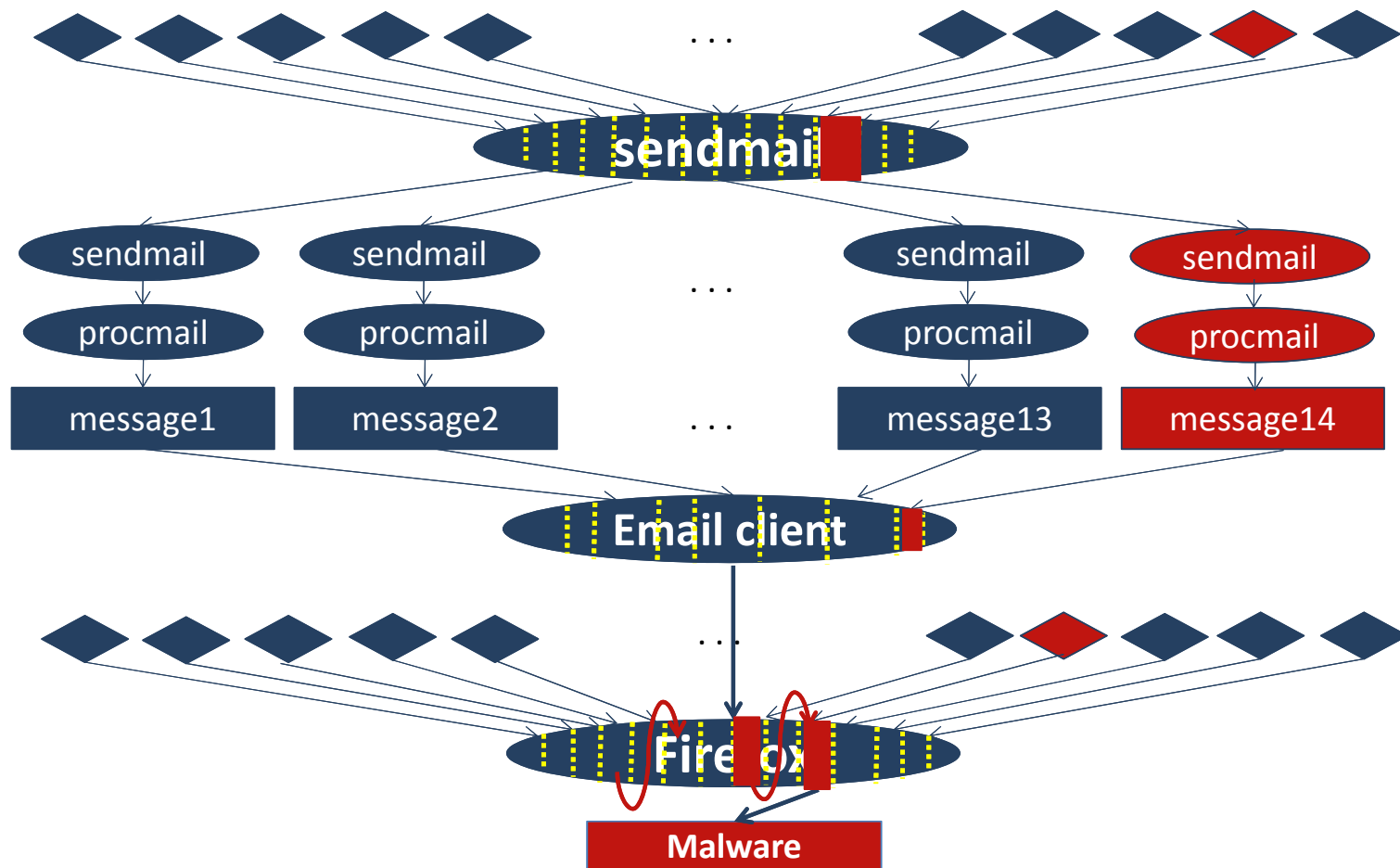
- Finer-grained subject : Execution “UNIT”
 - Dynamically partition the execution of a process into **autonomous** execution segments
 - Units are not always independent
 - Detect causality between units

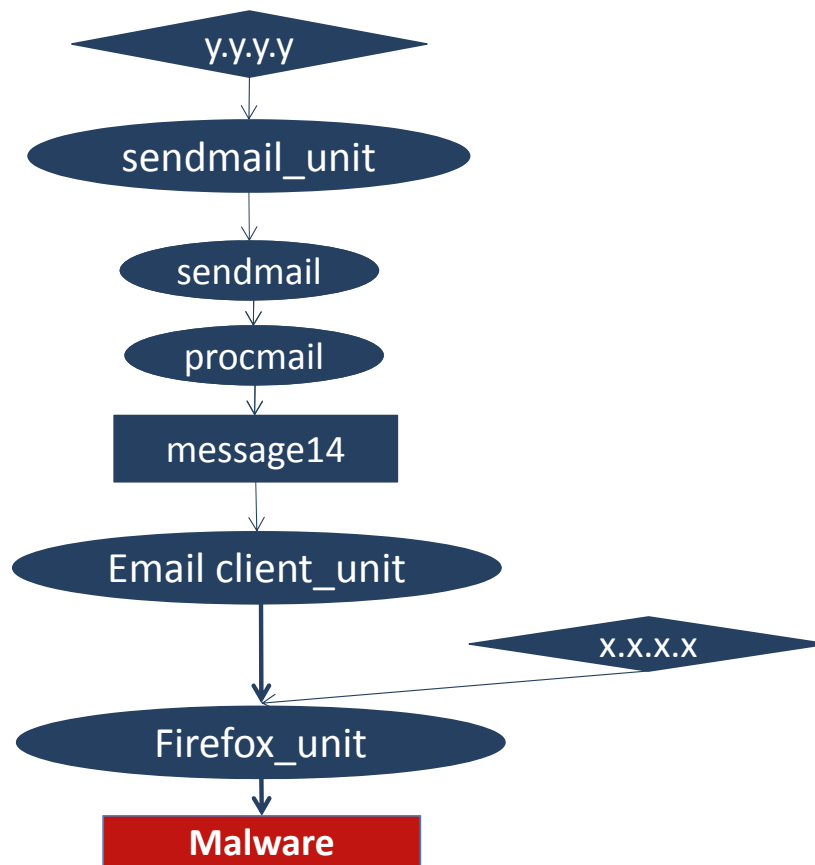


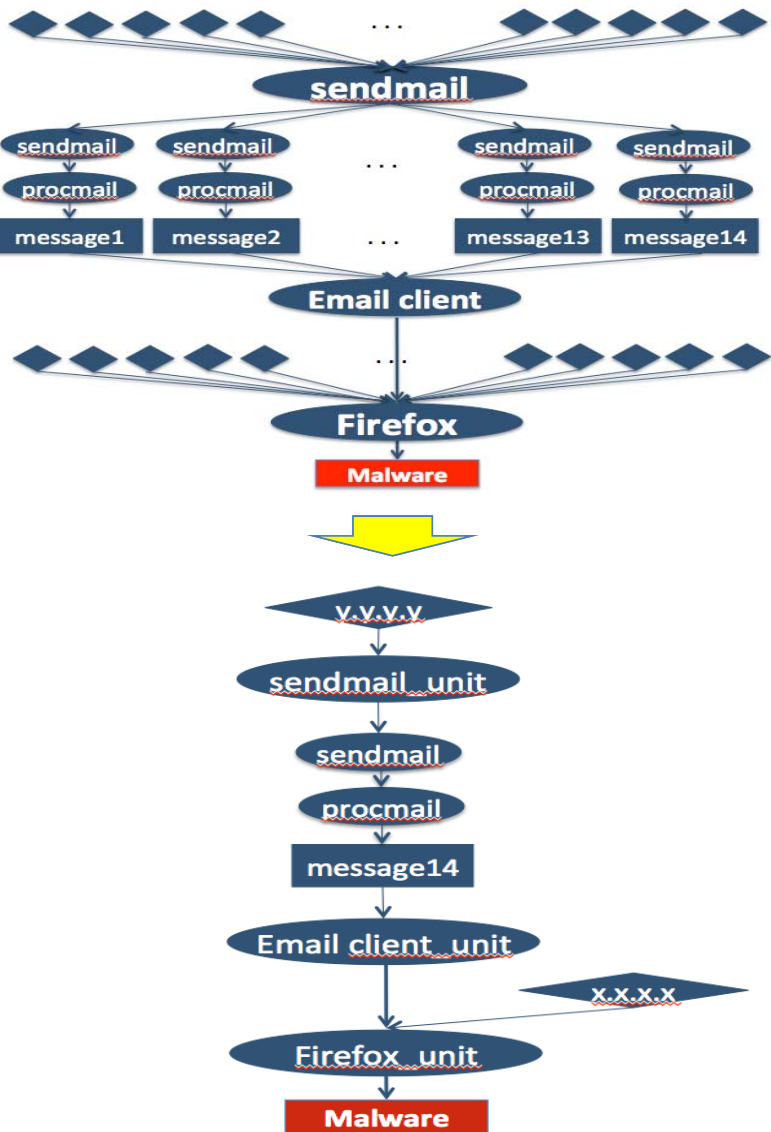
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 - Dynamically partition the execution of a process into **autonomous** execution segments
 - Units are not always independent
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Previous approaches [SOSP'03, SOSP'05] :

**51 Processes, 15 Files,
251 Network addresses, 351 Edges**

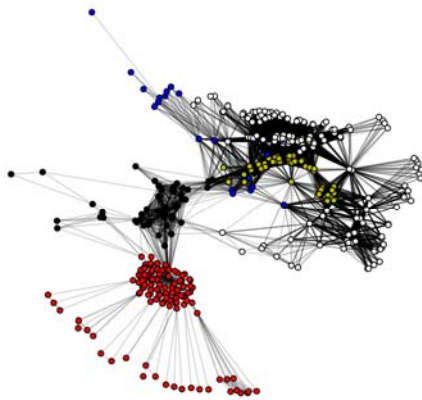
16.3 times smaller

BEEP :

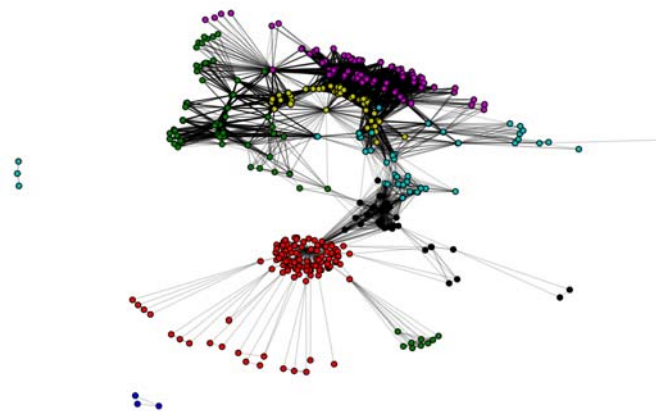
**10 Processes, 2 Files,
6 Network addresses, 23 Edges**

Evaluation on Real APT Attacks

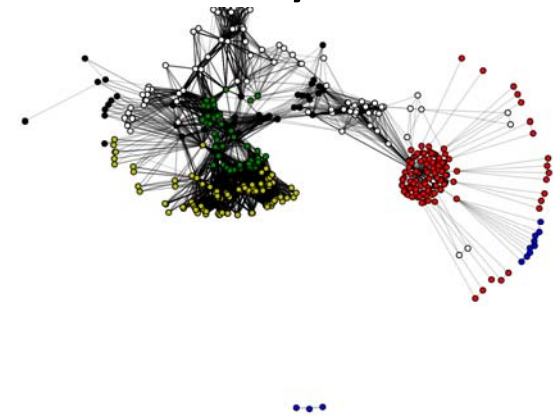
- Reproduce about 15 real attacks
- Our graphs are usually one-two orders of magnitude smaller, and cover whole attack paths with over 90% accuracy.



“Russian Campaign”



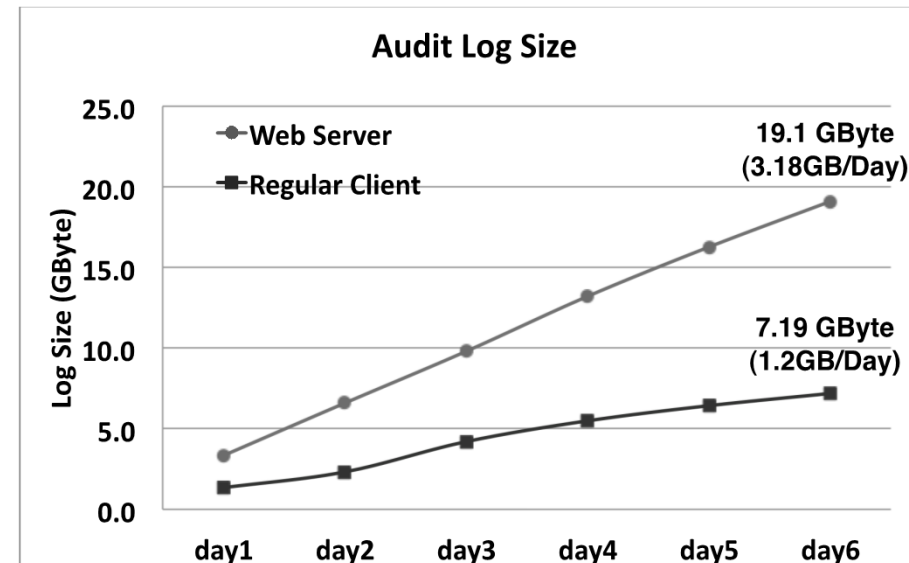
“Black Vine”



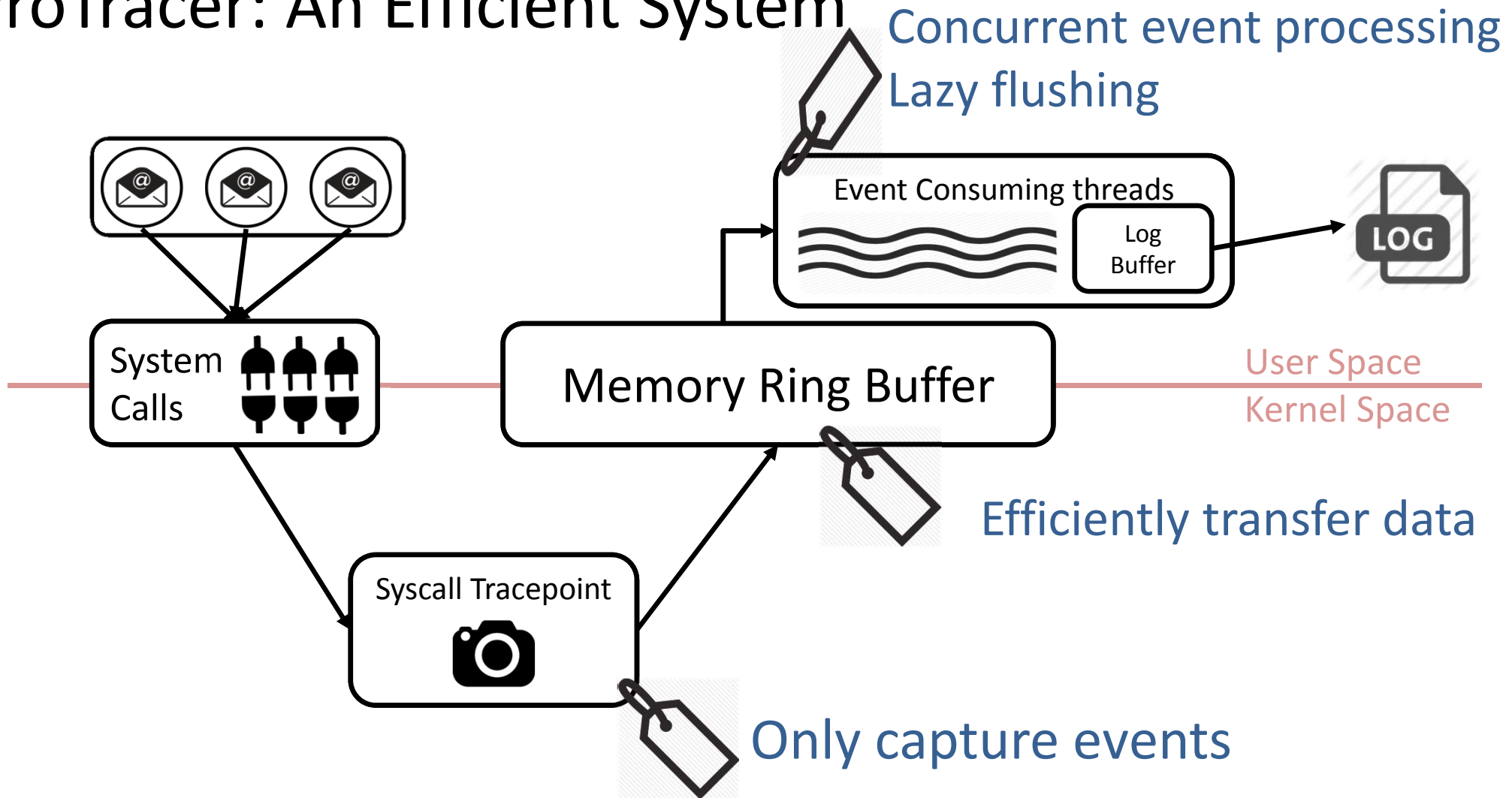
“Op-DeputyDog”

Limitations of Traditional Audit Logging

- Dependence explosion problem (**NDSS'13**)
 - Long running programs generate a lot of system level dependences during their lifetime
- High overhead
 - Linux Audit Framework: **~40%** run time slow down
 - Space overhead
 - Our solution
 - ProTracer (**NDSS'16 Distinguished Paper**)



ProTracer: An Efficient System

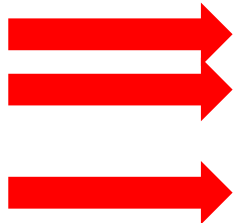


Example: Avoid *Redundant* Events

```

1. # vim opening a large file
2.   ...
3.   while((size = read(fd, buf)) > 0):
4.       add_node(root, buf)
5.       ...
6.   exit();

```



Logging

```

...
PID = 1483, TYPE = SYSCALL: Syscall = read
PID = 1483, TYPE = SYSCALL: Syscall = read
PID = 1483, TYPE = SYSCALL: Syscall = read
PID = 1483, TYPE = SYSCALL: Syscall = read
PID = 1483, TYPE = SYSCALL: Syscall = read
PID = 1483, TYPE = SYSCALL: Syscall = read
...
PID = 1483, TYPE = SYSCALL: Syscall = exit

```

ProTracer

```

...
T[ PID=1483 ] = { vim }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
T[ PID=1483 ] = T[ PID=1483 ] V { fd } = { vim, fd }
...
LogBuffer: T[ PID=1483 ] = { vim, fd }

```

Example: Lazy Flushing

1. # temporary files
2. `f = open(fname, create | write)`
3. # File manipulation on the file
4. `while (not done)`
5. `edit(f)`
6. # delete temporary file
7. `delete(f)`

Logging

```
...
TYPE = SYSCALL: Syscall = open, FD = 8
TYPE = SYSCALL: Syscall = write, FD = 8
.....
TYPE = SYSCALL: Syscall = write, FD = 8
TYPE = SYSCALL: Syscall = unlink, FD = 8
...
```

ProTracer

```
...
T[ FD=8 ] = { }
T[ FD=8 ] = { vim }
LogBuffer: T[ FD=8 ] = { vim }
T[ FD=8 ] = T[ FD=8 ] V { vim } = { vim }
LogBuffer: T[ FD=8 ] = { vim }
DEL: T[ FD=8 ]
...
```

LogBuffer

```
T[ FD=8 ] = { vim }
T[ FD=8 ] = { vim }
```



Evaluation: Storage Efficiency (3 months, client)

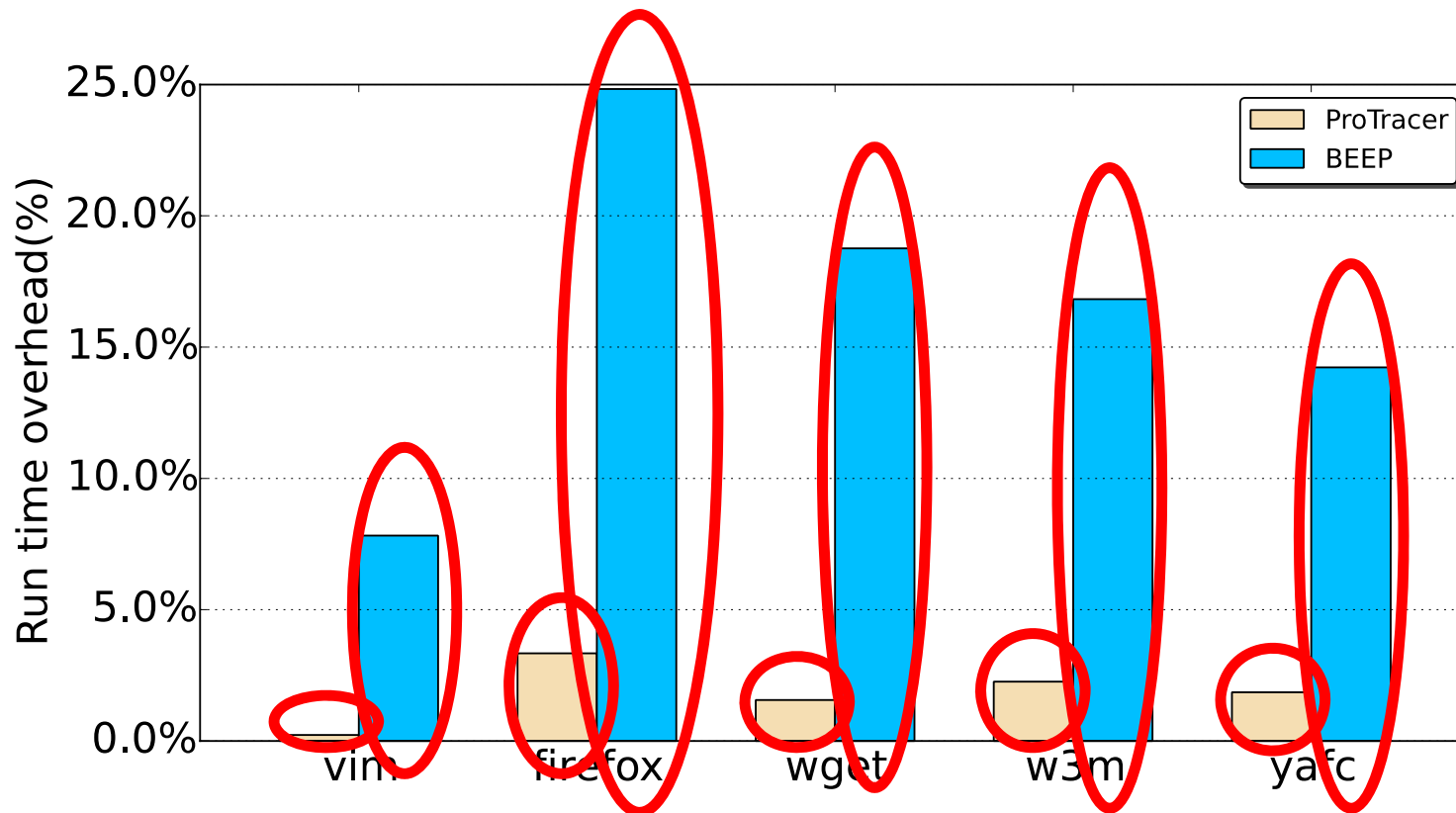
BEEP

[NDSS'13]

168,269,688 KB

ProTracer
2,437,010 KB

Evaluation: Run time Efficiency (Client Programs)



1.9%
v.s.
16.5%

Whole system: 7% v.s. 40%

Outline

Temporal Dynamic Analysis:

analyze program execution history

1. Attack investigation

2. Application vetting

1. Execution partitioning overcomes the dependence explosion problem
2. ProTracer reduces space consumption by from over 1GB per day to less than 20MB per day, and runtime overhead from 40% to 7%

on state

Dynamic Analyses

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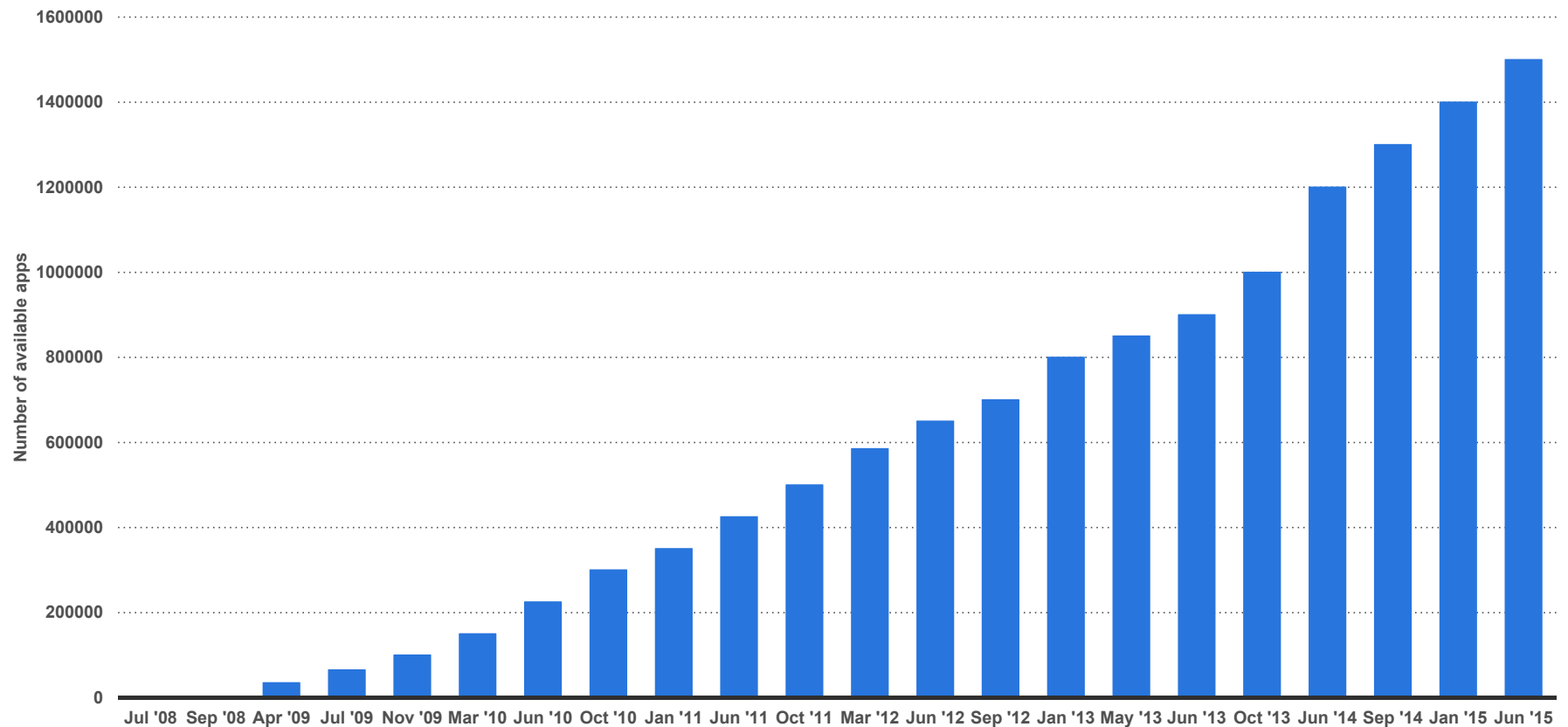
Spatial Dynamic Analysis:

analyze a snapshot of program execution state

3. Memory forensics

Dynamic Analyses

Motivation: Total number of iOS applications (2008-2015)



Motivation: ...but Apple is watching out for you

Apple has a pre-release app vetting procedure call **App Review**

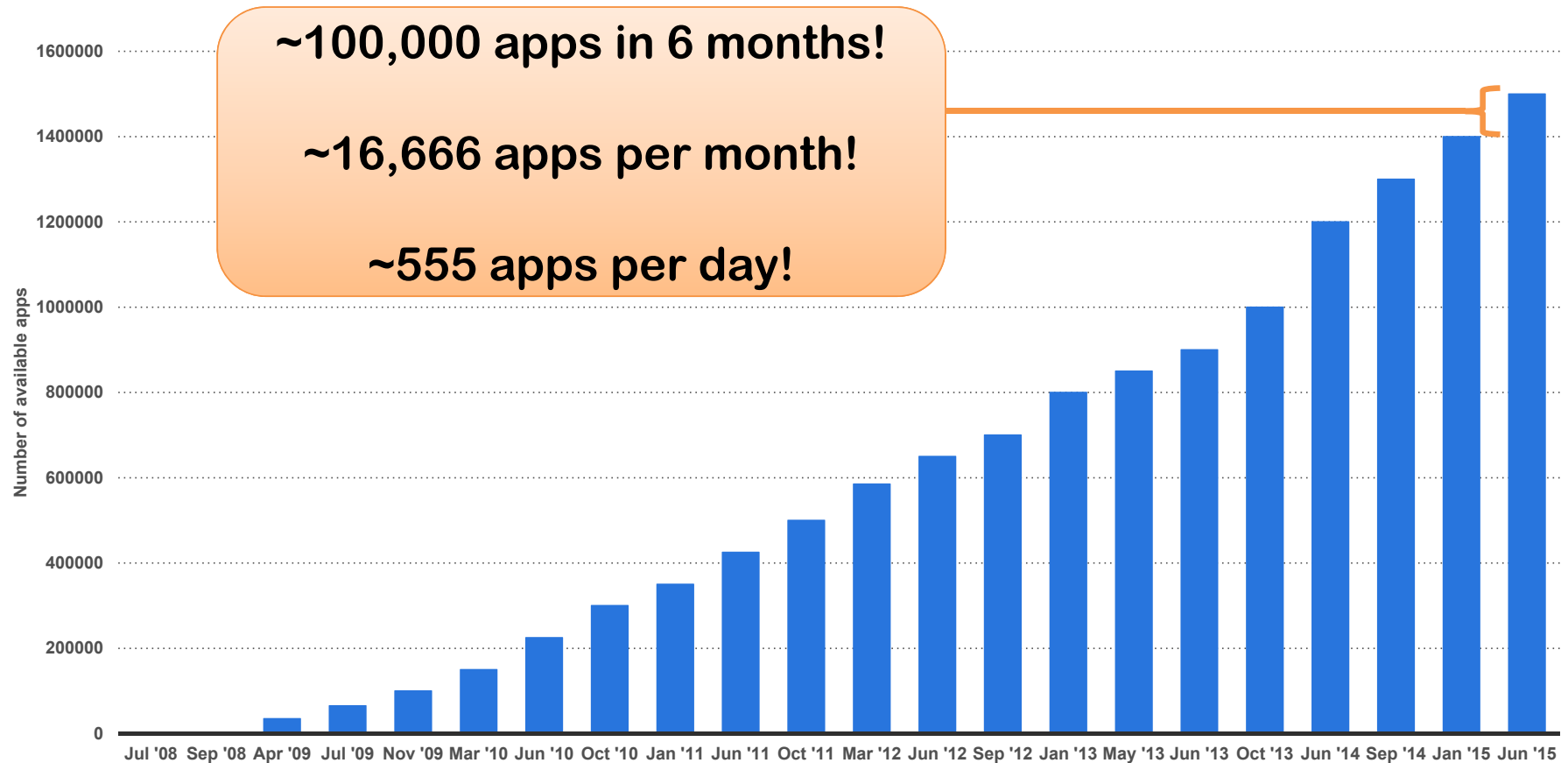
“We review all apps [...] to ensure they are reliable, perform as expected, and are free of offensive material.”

- Apple App Review Team

<https://developer.apple.com/app-store/review/>

Apple is secretive about the specific evaluation ... Security?

Motivation: Total number of iOS applications (2008-2015)



Observations...

App Review vetting procedure is:

- 1) Very fast and lightweight
- 2) “Almost” fully automated
- 3) Static Program Analysis

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App Review vetting procedure is:

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Easily Subverted!

Control Flow & Data Obfuscation

Wang et.al [USENIX Security'13]

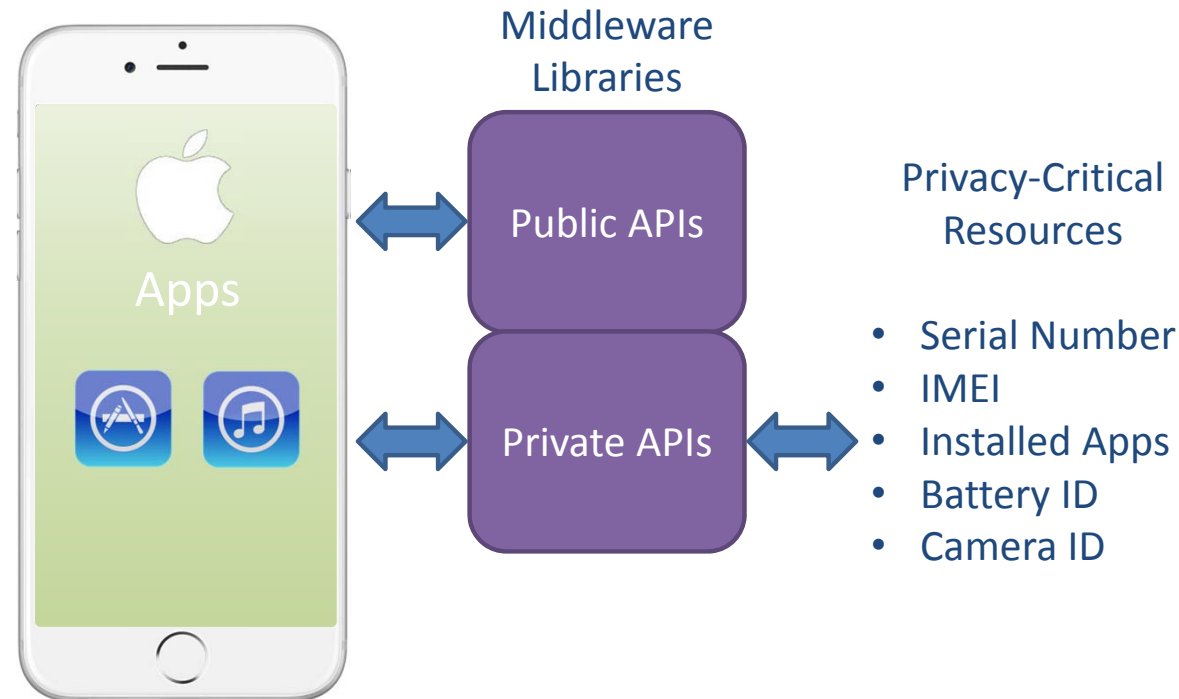
Han et.al [ACNS'13]

Zheng et.al [ASIACCS'15]

InstaStock, FindAndCall,
PawnStorm.A, LBTM, FakeTor,
XCodeGhost, etc....

Goal: Private APIs

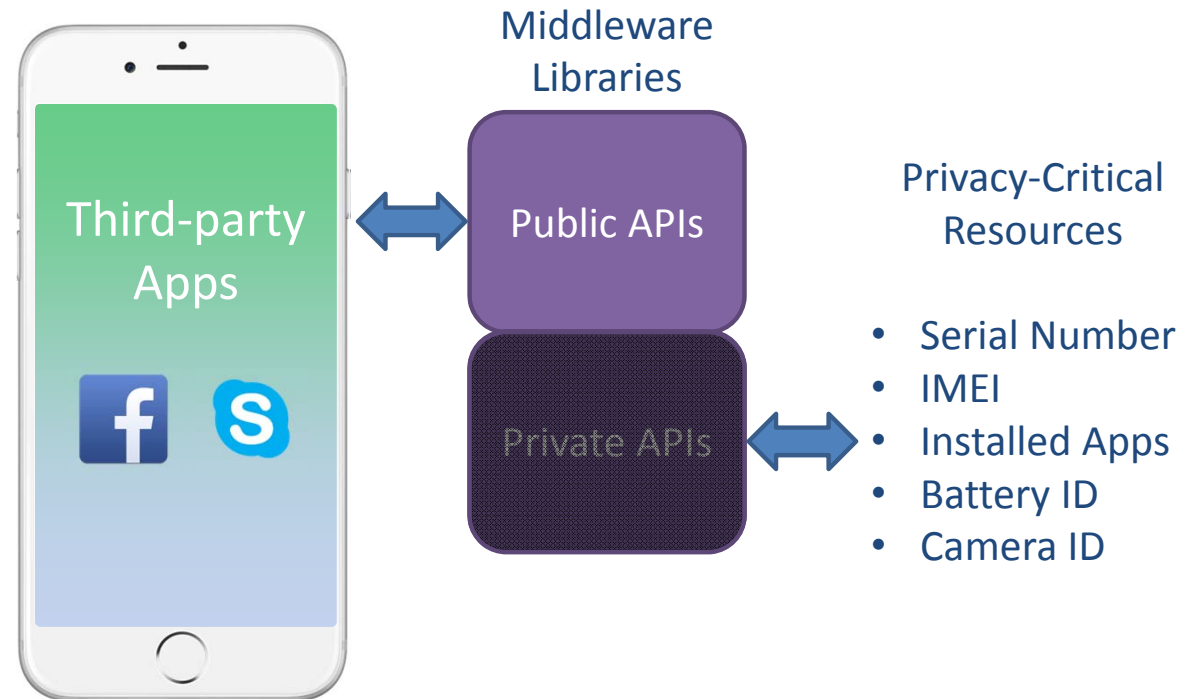
Middleware APIs reserved
only for Apple's internal use



Goal: Private APIs

Middleware APIs reserved only for Apple's internal use

Private APIs are not exported for third-party code to use

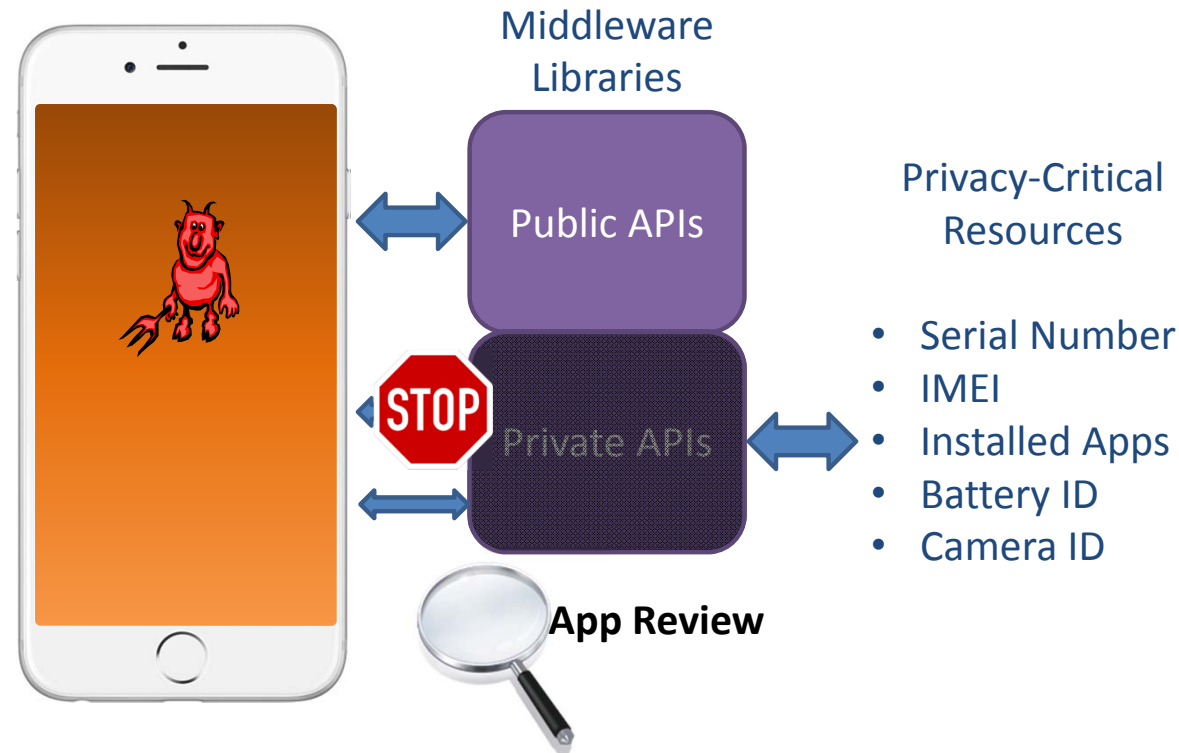


Goal: Private APIs

Middleware APIs reserved only for Apple's internal use

Private APIs are not exported for third-party code to use

Malware are still finding and invoking the Private API despite vetting



Private API Abuse is Difficult To Detect

iOS applications are mostly written in Objective-C &
Objective-C functions can be called sending a message to object

Invoke method **foo** on object **obj** with parameter **param**:

```
objc_msgSend ( obj, "foo", param );
```

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Message Selector = String

Private API Abuse is Difficult To Detect

iOS applications
Objective-C func

**Static Analysis will not be
able to detect complex
private API invocations!**

C &
message to object

Invoke method **foo** on object **obj** with parameter **param**:

```
char sel[3]; strcpy ( sel, "f" ); strcat ( sel, "oo" );
objc_msgSend ( objc, sel, param );
```

+ Obfuscation
+ Encryption

iRiS - Automated Vetting of Private API Abuse

Forced Execution



Drives the app's execution to
unresolved API call sites using
forced execution

Forced Execution

- Dynamic analysis engine that forces a binary to execute
 - Provide no inputs or any environment setup
 - Use random inputs
 - Force branch outcomes at a few places

Force Execution: Conditional Message Selector

```
1. char hidden[] = "\x73\x68\x75\x74\x64\x6f\x77"  
2.           "\x6e\x44\x65\x76\x69\x67\x67\x67";  
3. getAppsOrShutdown( void * obj, int key )  
4. {  
5.     char * sel;  
6.     if( time() == DEC_21_2017 )  
7.     {  
8.         sel = XOR(hidden, key);  
9.     }  
10.    else  
11.    {  
12.        sel = "allInstalledApplications";  
13.    }  
14.    objc_msgSend( obj, sel );  
15. }
```

iRiS must resolve the
value of **sel** at this call

Force Execution: Conditional Message Selector

```
1. char hidden[] = "\x73\x68\x75\x74\x64\x6f\x77"  
2.           "\x6e\x44\x65\x76\x69\x67\xb7";  
3. getAppsOrShutdown( void * obj, int key )  
4. {  
5.     char * sel;  
6.     if( time() == DEC_21_2017 )  
7.     {  
8.         sel = XOR(hidden, key);  
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10.    else  
11.    {  
12.        sel = "allInstalledApplications";  
13.    }  
14.    objc_msgSend( obj, sel );  
15. }
```

Not Resolvable by Static
Analysis

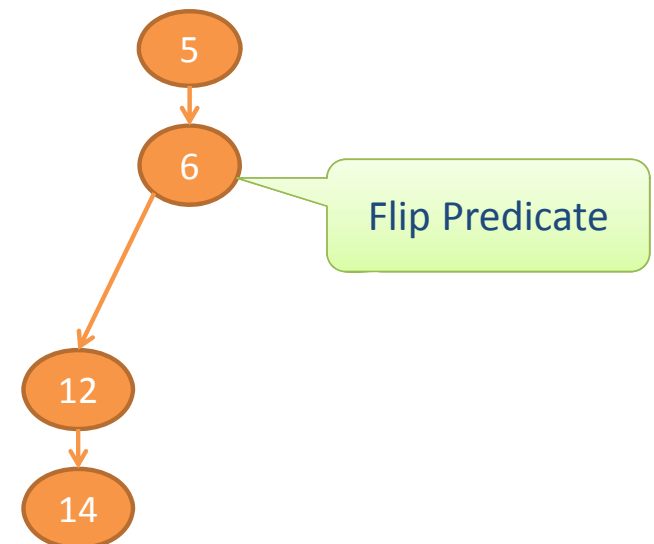
Force Execution: Conditional Message Selector

```

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2.                "\x6e\x44\x65\x76\x69\x67\x67";
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4. {
5.     char * sel;
6.     if( time() == DEC_21_2017)
7.     {
8.         sel = XOR(hidden, key);
9.     }
10.    else
11.    {
12.        sel = "allInstalledApplications";
13.    }
14.    objc_msgSend( obj, sel );
15. }

```

Most Executions follow
the false branch

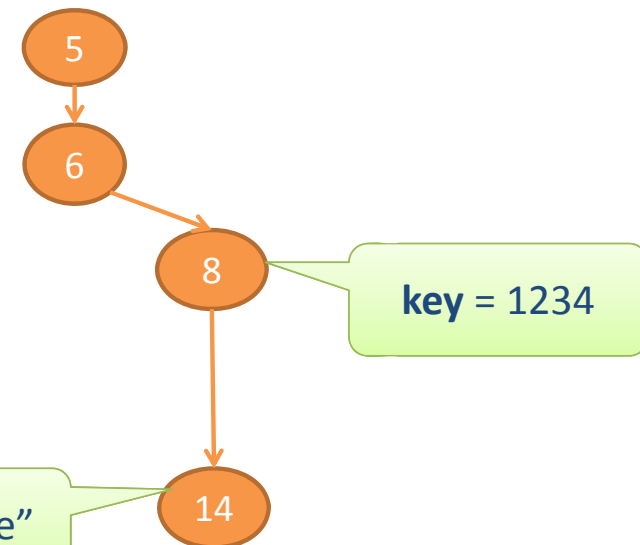


iRiS determines if a new call target can be resolved by flipping a predicate

Force Execution: Conditional Message Selector

```
1. char hidden[] = "\x73\x68\x75\x74\x64\x6f\x77"  
2.           "\x6e\x44\x65\x76\x69\x67\x67";  
3. getAppsOrShutdown( void * obj, int key )  
4. {  
5.     char * sel;  
6.     if( time() == DEC_21_2017 )  
7.     {  
8.         sel = XOR(hidden, key);  
9.     }  
10.    else  
11.    {  
12.        sel = "allInstalledApplications";  
13.    }  
14.    objc_msgSend( obj, sel );  
15. }
```

Force Execution
Resolved Code Path



Crash-free Execution Model

- Ideas on memory access exception

- Skip it?

- A lot of following exceptions, cascading effect on program state corruption
 - Lose heap data

- Allocate a piece of memory on demand

- It is not sufficient by just fixing the corrupted pointer itself
 - Fix the other correlated pointers

```
1. if (time() == DEC_21_2017)
```

```
2.   p=0;
```

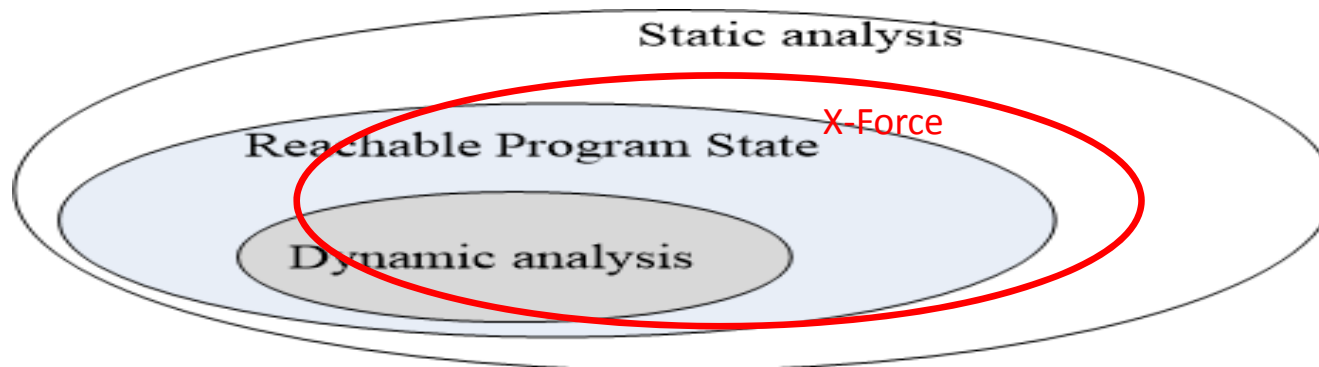
```
3.   q=p+5;
```

```
4.   *q=...
```

Flipped to
true

Crash on *q

The Essence of X-Force



- X-Force is neither sound nor complete
 - Unsound due to path infeasibility and violation of input precondition
 - Incomplete due to the prohibitively large search space
- A pragmatic solution for certain applications
 - Fast
 - Path feasibility may be violated at only a small number of places
 - Less than 10 dynamic predicates out of hundreds of thousands.
 - Naturally handle packed, obfuscated, and even self-modifying binaries
 - Existing dynamic analysis can be easily ported to X-Force
 - Good at exposing behavior, cannot generate exploit inputs

iRiS Evaluation

Vetting platform: iPad 4 + Our ported Valgrind (iOS)

2019 free apps from official App Store

- 9 categories

- Popular apps listed in iTunes Preview

- Crawled in March 2015

- App size: 1-80MB, median 3MB

iRiS Evaluation

149 (7%) applications use private APIs

Identified a total number of 153 private APIs

Apple brought down most of these apps from their store after our paper was published

<http://www.digitaljournal.com/technology/apple-blacklists-hundreds-of-apps-that-stole-personal-user-data/article/446995>

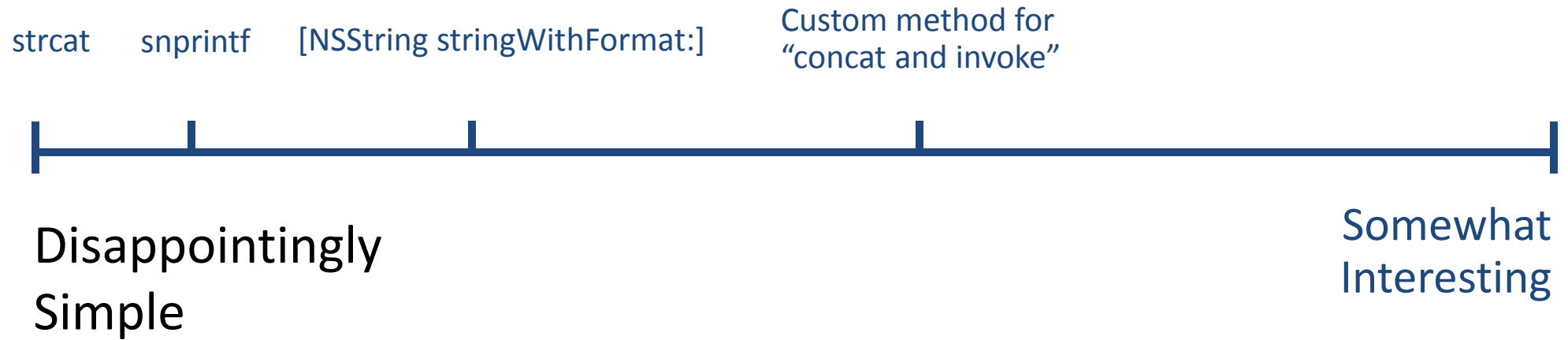
Private APIs: Access User Identification

Framework	API Name	Functionality	#Apps
SpringBoardServices	SBSSpringBoardServerPort	Initialize port with Springboard	3
	SBSCopyApplicationDisplayIdentifiers	Obtain ids of all running apps	3
	SBFrontmostApplicationDisplayIdentifier	Obtain id of the front most app	3
	SBSCopyLocalizedApplicationNameForDisplayIdentifier	Get app name from its id	33
MobileCoreServices	[LSApplicationWorkspace defaultWorkspace]	Obtain the default workspace object	2
	[LSApplicationWorkspace allApplications]	Get all installed apps	1
	[LSApplicationWorkspace allInstalledApplications]	Get all installed apps	1
	[LSApplicationWorkspace applicationIsInstalled:]	Check if a specific app is installed	1

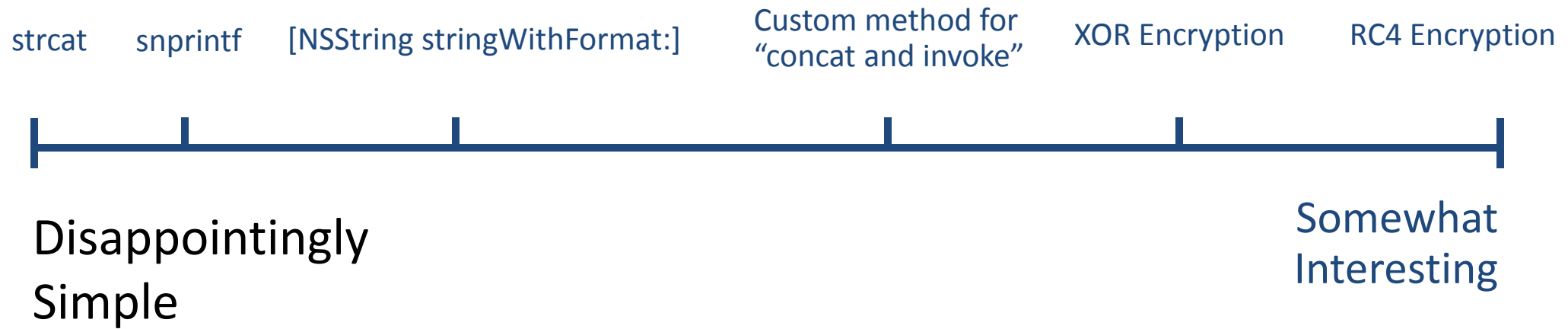
Private APIs: More User Identification

Framework	API Name	Functionality	#Apps
AppleAccount	[AADeviceInfo appleIDClientIdentifier]	Obtain the Apple ID of the device user	1
AdSupport	[ASIdentifierManager sharedManager]	Obtain a reference to the ASID manager	25
	[ASIdentifierManager advertisingIdentifier]	Obtain the device's ASID	25
	[ASIdentifierManager advertisingTrackingEnabled]	Check if advertising tracking is enabled	23
IOKit	IOMasterPort	Initialize communication with IOKit	21
	IOServiceMatching	Find and open the specified IOService object	21
	IOServiceGetMatchingService		21
	IORegistryEntryCreateCFProperty	Locate the specified property (e.g. S/N)	19

Range of Obfuscation Tactics



Range of Obfuscation Tactics



Outline

Temporal Dynamic Analysis:

analyze program execution history

1. Attack Investigation

2. Application vetting

Spatial Dynamic Analysis:

analyze a snapshot of program execution state

3. Memory forensics

Dynamic Analyses

Outline

Temporal Dynamic Analysis:

analyze program execution history

1. Audit logging

2. Forced execution

Spatial Dynamic Analysis:

analyze a snapshot of program execution state

3. Memory forensics

Dynamic Analyses

Execution State (in Memory) Contains Rich Information

Running Processes

Volatile IPC Data

Executing Malware

Decrypted Private Data

Open Network Connections

Encryption Keys

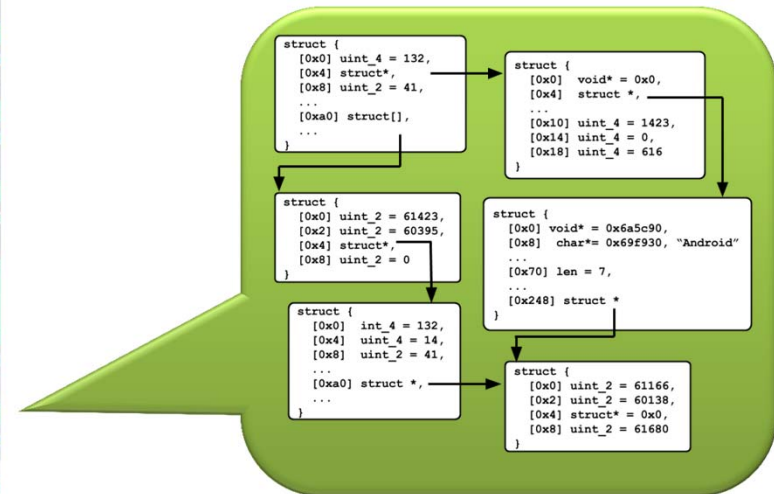
Application Data
(Browser History, Chat Logs, ...)

Much More ...



In-Memory Information Reverse Engineering

Information in memory
is stored in Data
Structures



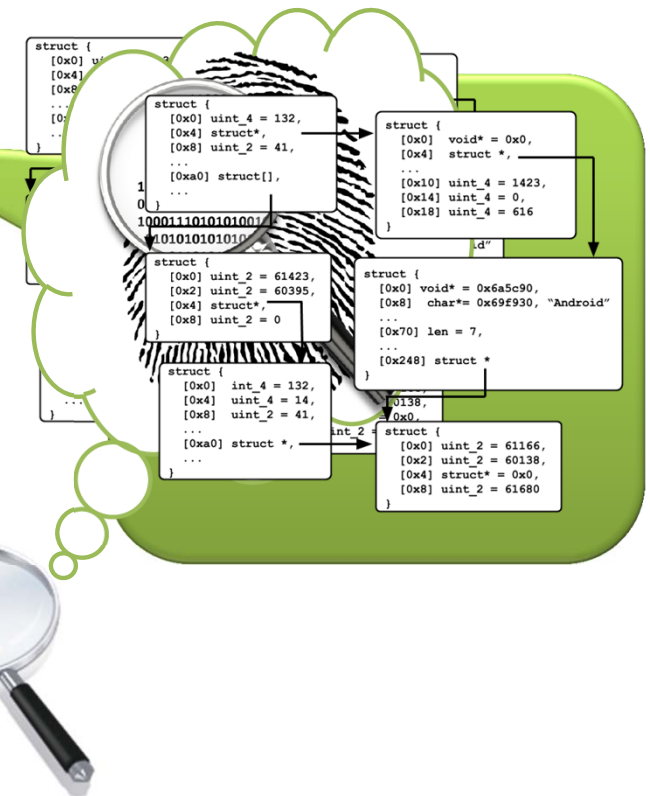
In-Memory Information Reverse Engineering

Investigators collect
a memory image

Build signatures of
data structures to
recover

Scan the memory
image to find data
structure instances

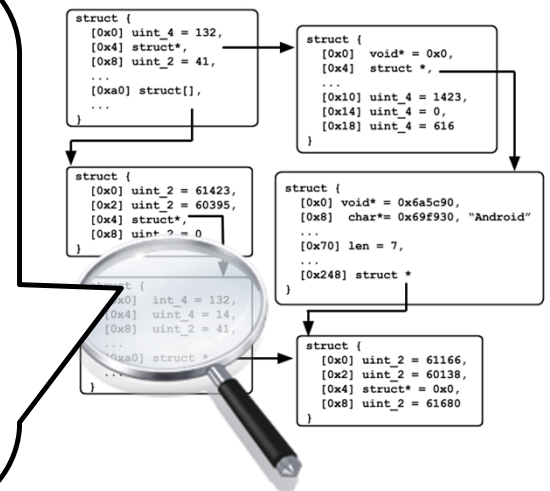
2FF0	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
3000	AF 21 F4 31 CD 1F 31 21 EB 31 CD 1F 31 CD 89 46
3010	21 00 50 11 F0 03 CD 8F 46 21 00 58 11 FF 07 36
3020	70 23 1B 7B B2 20 F8 CD 98 46 CD B9 46 CD 3E 00
3030	CD 1B 00 FE 53 CA 3F 30 FE 31 CA 00 00 18 F1 06
3040	BE AF 21 2C 31 77 23 10 FC 3E 18 32 FE 31 21 50
3050	D0 11 98 03 CD 8F 46 AF 32 26 31 32 4F 31 21 3E
3060	31 CD 1F 31 32 F7 31 32 F8 31 3E 03 32 25 31 21
3070	CB 52 22 27 31 3E 08 32 F9 31 CD 8D 47 0E 16 AF
3080	32 DE 31 CD 52 41 CD 95 33 3E 01 32 2B 31 3E 08
3090	32 D9 31 AF 32 DB 48 32 DC 48 32 DF 48 32 DE 48
30A0	32 DD 48 21 00 00 22 4A 31 22 4D 31 21 00 01 22
30B0	54 31 11 56 31 CD A8 35 11 6F 31 CD A8 35 11 88
30C0	31 CD A8 35 11 A1 31 CD A8 35 11 BA 31 CD A8 35
30D0	11 00 00 AF CD 33 00 CD 3E 00 21 56 31 11 17 00
30E0	19 36 20 11 19 00 19 36 30 19 36 50 19 36 60 19
30F0	36 20 CD BF 48 C3 01 32 06 13 21 63 D0 11 27 00
3100	3E 76 77 19 10 FC 06 13 21 64 D0 11 29 00 3E 77
3110	77 19 10 FC 06 28 0E 3C 21 28 D0 CD 56 3D C9 77
3120	23 77 23 77 C9 03 00 BA D2 00 00 01 00 01 01 03
3130	00 8B 50 02 0D 00 1F 3C 0A 04 3E 00 00 20 3C 31
3140	3C 8B 53 20 00 00 00 17 54 00 0E 00 01 00 00 03
3150	FF D0 00 00 00 01 00 00 00 00 00 00 00 00 00 00
3160	00 00 00 00 00 00 00 00 00 00 00 00 00 20 00 C0
3170	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 C0



(Prior) State of the Art: Data Structure Recovery

Evidence is
recovered from
plain-text or self-
contained fields

```
struct user_account {
[0x00]  short int u_type;
[0x04]  pid_t u_pid;
[0x08]  char u_line[32];
[0x28]  char uid[4];
[0x2C]  char user[32];
[0x4C]  char password[128];
[0xCC]  char u_host[128];
[0x14C] short e_termination;
...
}
```



A Cyber-Crime

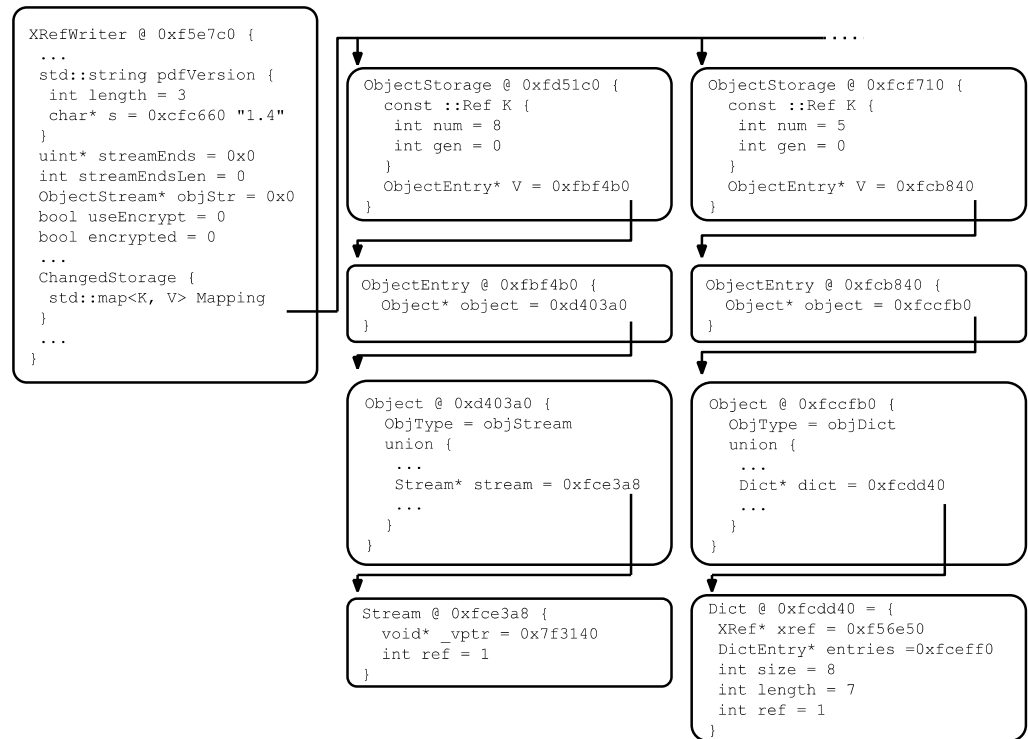
Based on true events that occurred
at the authors' university...

State of the Art ... but Limited

Finds raw data structure instances in memory image

Still cannot understand the **content** of the data structure!

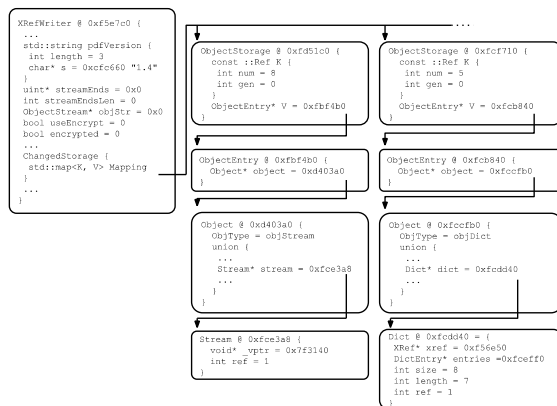
E.g., images, documents, formatted/encoded data



DSCRETE: Content Reverse Engineering

[Usenix Security Best Student Paper]

The application that defined the data structure contains printing/rendering logic for it too!



Official Grade Report
Official University Document

Course: CS123 Room: CSB 401 ID: 90123

Name	Homework	Exam	Final
John Taylor	A	C	A
Jason Zhang	A	B	A
Julian Smith	A	C	B
Christopher Zhu	F	C	D
Brett Choi	D	A	B
Sandy Chen	B	C	C
Sunal Vanna	A	A	A

DS-CRETE Content Reverse Engineering

[Usenix Security Best Student Paper]

Input: Data Structure Instance

Output: Formatted Content

Printing/Rendering Logic

Program Code

```
struct pdf* my_pdf;
my_pdf = load_pdf_file(...);
main_loop(my_pdf); // User edits PDF
save_pdf_file(my_pdf);
exit(0);
```

Rendering Function


```
save_pdf_file(struct pdf* ptr)
{
    char* buf = format_pdf(ptr);
    fwrite(buf, ...);
}
```

Scanner+Renderer

```
110100000101010  
111101001011100  
011010010100101  
001000100100111
```

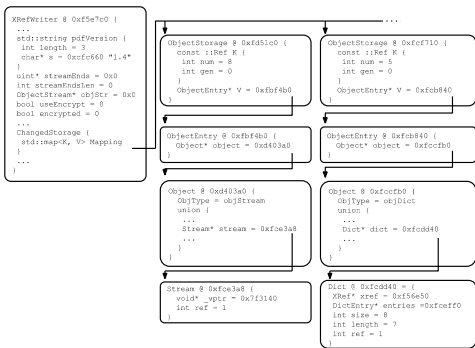
Rendering Function

```
save_pdf_file(struct pdf* ptr)  
{  
    char* buf = format_pdf(ptr);  
    fwrite(buf, ...);  
}
```



Intuition: Invalid input will crash the rendering logic

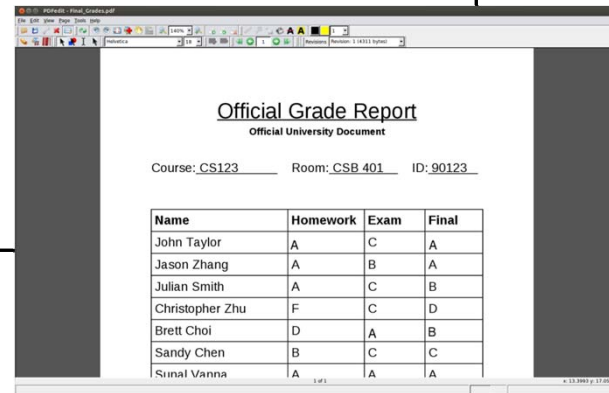
Scanner+Renderer



Rendering Function

```
save_pdf_file(struct pdf* ptr)
```

```
{
  char* buf =
  fwrite(buf,
}
```



Present every offset of a memory image to the rendering logic &
Reported the valid output

Cross-State Execution

Identified Candidate

Program Code

```
struct pdf* my_pdf;  
my_pdf = load_pdf_file(...);  
main_loop(my_pdf); // User edits PDF  
save_pdf_file(my_pdf);  
exit(0);  
  
save_pdf_file(struct pdf* ptr)  
{  
    char* buf = format(ptr);  
    fwrite(buf, ...);  
}
```

Cross-State Execution

App's Memory



Program Code

```

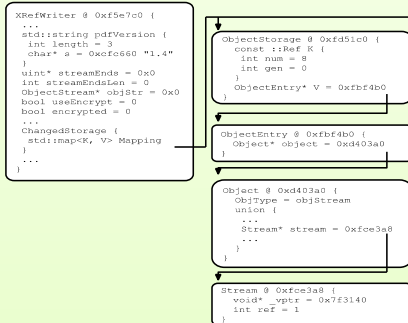
struct pdf* my_pdf;
my_pdf = load_pdf_file(...);
main_loop(my_pdf); // User edits PDF
save_pdf_file(my_pdf);

exit(0);

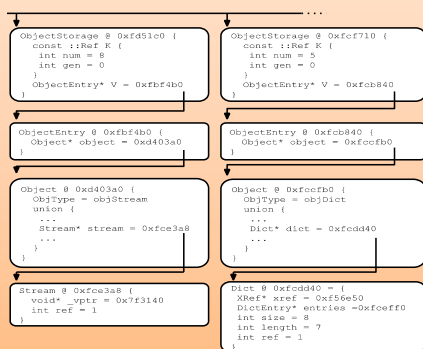
save_pdf_file(struct pdf* ptr)
{
  char* buf = format(ptr);
  fwrite(buf, ...);
}
  
```

Cross-State Execution

App's Memory



Memory Snapshot



Program Code

```
struct pdf* my_pdf;
my_pdf = load_pdf_file(...);
main_loop(my_pdf); // User edits PDF
save_pdf_file(my_pdf);
exit(0);
```

Begin Cross-State Execution!

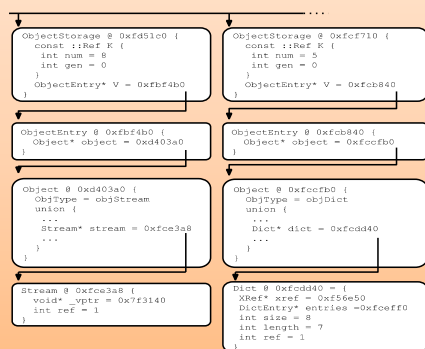
1. Map in memory snapshot
2. Swap my_pdf pointer

Cross-State Execution

App's Memory



Memory Snapshot



Program Code

```

struct pdf* my_pdf;
my_pdf = load_pdf_file(...);
main_loop(my_pdf); // User edits PDF
save_pdf_file(my_pdf);
exit(0);
  
```

```

save_pdf_file(struct pdf* ptr)
{
  char* buf = format(ptr);
  fwrite(buf, ...);
}
  
```

Let's catch that criminal...



Back at the Forensics Lab...

App.

convert

gnome-paint

gThumb

gnome-screen

PDFedit - Final_Grades.pdf

File Edit View Page Tools Help

140%

Revisions: Revision: 1 (4311 Bytes)

Xfig - usa.fig

File Edit View Snap Help usa.fig

1 object(s) saved in "usa.fig"

Drawing

Editing

Zoom 1 Grid Mode Point Posn

Report

ment

401 ID: 90123

Exam	Final
C	A
B	A
C	B
C	D
A	B
C	C
A	A

0.26, 0.16, 0.10
1 zombie
0 hi, 0.0 si, 0.0 st
255068 buffers
14369524 cached

TIME+ COMMAND

32.62 compiz
10.37 Xorg
04.60 firefox
10.80 kworker/2:1
01.60 init
00.02 kthreadd
04.18 ksoftirqd/0
00.02 migration/0
00.97 watchdog/0
00.04 migration/1
06.06 kworker/1:0
03.14 ksoftirqd/1
00.84 watchdog/1
00.02 migration/2

libreoffice-writer.png - 8/...

Previous Next

128.10.125.22 - - [23/Jul/2013:16:36:10] "GET / HTTP/1.1" 304 0 "-"
"Mozilla/5.0 (X11; Linux x86_64) AppleWebKit/537.36 (KHTML, like Gecko)
Chrome/28.0.1500.71 Safari/537.36".

128.10.125.21 - - [23/Jul/2013:16:42:25] "GET / HTTP/1.1" 304 0 "-"
"Mozilla/5.0 (X11; Linux x86_64) AppleWebKit/537.36 (KHTML, like Gecko)
Chrome/28.0.1500.71 Safari/537.36".

128.10.125.22 - - [23/Jul/2013:16:44:07] "GET / HTTP/1.1" 304 0 "-"
"Mozilla/5.0 (X11; Linux x86_64) AppleWebKit/537.36 (KHTML, like Gecko)
Chrome/28.0.1500.71 Safari/537.36".

8 files (2.7 kB) · 1 file selected (26... 300% 32 x 32 · 267 bytes · 10/08/2012 03:36:45 PM

SQL log

788

7

VCR: Reverse Engineer Preview Pictures

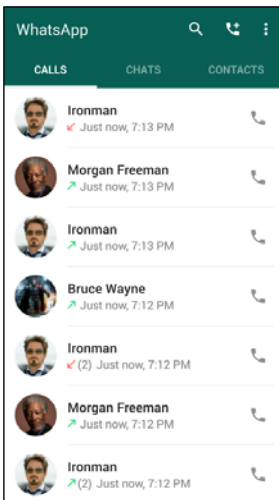


Photographic Evidence is available as soon as a camera app is opened, even without taking a photo or video

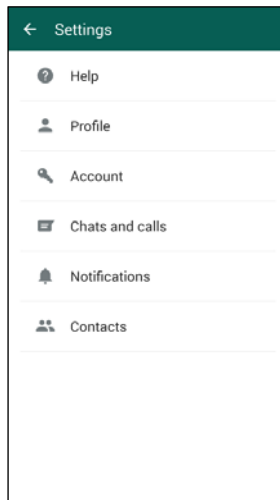


Camera preview is always buffering frames

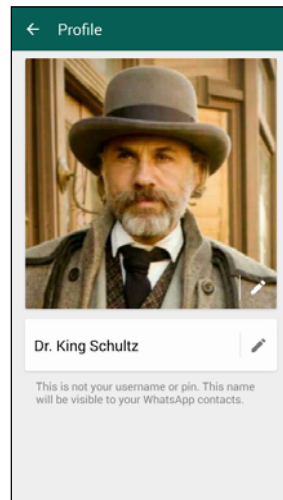
GUI: Reverse Engineer User Interfaces



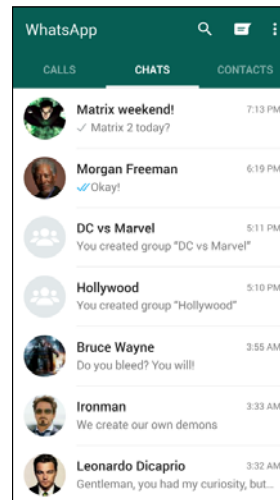
Screen -6



Screen -5



Screen -4



Screen -3



Screen -2



Screen -1



Screen 0

Outline

Temporal Dynamic Analysis:

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Dynamic Analyses

Outline

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3. Memory forensics

Dynamic Analyses

Thank you!
Questions?